

CONCLUSION BIBLIOGRAPHY

APPENDICES

CONCLUSION

One of the main goals of the project was to gain a good experience in compiler methodology. This goal has been achieved ; we also developed principles and techniques in the run-time system design as well as in the static management and in their interface.

By implementing the language in its whole, we had to control a huge bulk of information and one may ask oneself if this is really worthwhile. To this question we answer that complexity is a problem in itself ; having succeeded to master it in a reliable way is quite an achievement ; this has been made possible by carefully choosing basic principles and applying them in a systematic and modular way.

Let us now try to evaluate the translation process on the basis of the following criteria ; *efficiency, security, portability* and *design effort*.

Beforehand, a number of general considerations must be made.

On the one hand, some of these criteria are many-sided, e.g. the efficiency must be split up into compile-time and run-time efficiency and both of these have two aspects : time and space. On the other hand, these criteria are conflicting :

- an increase in efficiency, security and portability must be paid by an increase in design effort.
- a higher level of security and portability generally decreases the efficiency.
- a higher level of run-time efficiency generally decreases the compile-time efficiency in space and/or time.
- a higher run-time efficiency in space generally causes a decrease in run-time efficiency in time and vice versa.

In the implementation described in this book the stress has been laid on run-time efficiency, security and portability while keeping a reasonable degree of compile-time efficiency. In a number of cases, some compromises have been made in order to keep the design effort inside reasonable limits. However the principle according to which "run-time efficiency of simple language features must not be affected by the presence of more intricate features" (Samelson-Bauer) has been constantly cared for.

Now, the evaluation proper can be formulated, it is based on the actual implementation on the X8.

- (1) *Compile-time efficiency in time* fits into quite acceptable limits ; it amounts to

$$5 + 6x \text{ seconds}$$

where x is the length of the source program in pages (one page = 60 average lines)

[4] .

- (2) *Compile-time efficiency in space* : the whole ALGOL 68 system including debugging facilities and initialized tables occupy 100K memory words of 27 bits, among which about 70K instructions. However the use of overlay techniques allows to run the compiler within a 32K words direct access memory.
- (3) *Run-time efficiency in time* : comparisons with the ALGOL 60-X8 compiler show an average increase in efficiency of 30% in favour of the ALGOL 68-X8 compiler. This is not negligible given the ALGOL 68 compiler has few restrictions, given its level of portability and the simplicity of the mechanism of local optimizations.
- (4) *Run-time efficiency in space* : this efficiency is difficult to estimate, we have no comparative figures at our disposal. Comparing the length of the object program with the length of the source program is meaningless : one single assignation ($x:=y$) may give rise to a variable number of objects instructions depending on the mode of x and y , i.e. on the data structure being assigned.
- (5) *Portability*

We distinguish two main aspects in portability :

- the language in which the compiler is written (a).
- the algorithm itself of compilation (b).

(a) The X8-compiler has been designed in an ALGOL 68-like language but hand-coded in assembly language. The language used for the design has been defined as our experience was growing ; as a consequence, the design needs some polishing before it can be accepted by a compiler.

(b) The second aspect, i.e. the portability of the algorithm itself of compilation (in particular the code generator) has been solved quite satisfactorily in the X8 compiler :

- the interface with the operating system is pretty well localized and can be easily modified.
- up to the production of intermediate code, only routines of lexical analysis dealing with internal representation of numbers (and possibly strings) have to be reconsidered to be transferred on a new hardware.
- the machine code generation itself, as it should be clear from this book, is surprisingly portable even on computers with a minimal hardware (III.5.1). For register allocation, only declarations making the correspondence between formal and actual registers must be specified according to each particular hardware.

Moreover, only a few routines have been made machine dependent in order to take direct advantage of particular hardware facilities.

These routines are :

<i>CONVERTACCESS</i>	(III.1.3)
<i>INREGPS</i>	(III.1.2)
<i>DEREFPS</i>	(III.1.2)
<i>GMI</i>	(III.2.2)
<i>COPYCELLS</i>	(III.5.4.2)

Also *GTAB*, (section III.2.2) containing instructions to be generated in an interpretative form, and its interpretation routine should be rewritten.

Finally, the global optimizer and the loader should probably be modified.

(6) *Design effort*

Roughly speaking, the compiler has consumed 20 men-years, but what does it really mean?

- The reader will be aware of how far we have optimized, and how few restrictions we have introduced.
- The programming tools we had at our disposal were very poor : the X8 assembler.
- The hardware and in particular the memory at our disposal was underdimensioned. As an example we had no protected backing store to save the compiler ; this means that it had to be reintroduced from cards and paper tape at each set of corrections or additions!
- The majority of the members of the team were unexperienced when they entered the project.
- In the 20 men years, the time spent for learning the language is incorporated, and we started reading early versions of drafts. Also, time spent for programming the special purpose system supporting the compiler and all debugging facilities is incorporated in the 20 men years.
- Finally, we must add that much time has been spent in checking the compiler carefully, step by step ; only a very small number of easily locatable bugs have been discovered since the compiler has been operational. It must be stressed that during the debugging process, the existence of high-level (design) and low-level (programming) documentation, carefully kept up to date, has appeared to be of the utmost importance.

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APPENDIX 1 : ANOTHER SOLUTION FOR CONTROLLING THE WOST% GARBAGE COLLECTION INFORMATION.

The idea of the solution is not ours, but it seems to be originated from the ALGOL 68-R compiler implemented on the ICL series by I. Currie. The solution is very efficient as such, however, it will be shown how, combined with our system, it would give still better results. It is to be noted that in such a combination practically the whole *gc* management described in this book remains valid.

The solution consists in constructing at compile-time a table representative of all possible *SWOST%* contents ; let us call it *GCTAB*. Each table element consists of the garbage collection information for one *WOST%* value (mode and access for example). Moreover elements are linked by a chain field in such a way a pointer to a table element gives access (through the chain field) to all the elements representative of a *WOST%* contents at a given moment.

Suppose for example the *WOST%* contents varies as follows :

```
(1) A
(2) A B
(3) A B C
(4) A B
(5) A
(6) A D
(7) A D E
```

The corresponding chain and entry points are sketched like this :

```
(1), (5)  A  ┌──┐
(2), (4)  B  ┌──┐
(3)       C  ┌──┐
(6)       D  ┌──┐
(7)       E  ┌──┐
```

With this table, instead of having a dynamic *GCWOST%*, the garbage collection information for each *BLOCK%* reduces to a pointer *gcw%* stored in its *H%*, pointer to a *GCTAB* element. In principle, instructions are generated to update the *gcw%* of the current *BLOCK%* each time a value is stored or deleted from *WOST%* ; the corresponding static management requires a field on *BOST (gc1)* representative of the *GCTAB* entry point associated with the value.

Two optimizations are now possible :

- (1) The first one corresponds to remark 2 at the end of I.2.4.3 : *gcw%* must not be updated at run-time if no garbage collection may take place during the time the corresponding value stays on *WOST%*. We can easily keep track of this fact by means of two new fields on *BLOCKTAB*, *gcp* and *gci* which are both pointers to *GCTAB*.

-*gcp* represents at each moment of the generation the actual *WOST%* state of the *BLOCK%*.

$-gci$ is a static image of $gcw\%$ of that $BLOCK\%$. Code updating $gcw\%$ will only be produced :

- (a) when code risking to activate the garbage collector is produced, and
 - (b) when gcp and gci of the corresponding block are different.
- (2) The second one corresponds to the minimization of the garbage collection information as explained in I.2.4.3.b, at the exception that what is minimized is the contents of $GCTAB$ instead of $GCWOST\%$. Clearly, the optimization is less advantageous here ; however some run-time actions may be saved given gcp will have to change less often and hence, instructions updating $gcw\%$ will have to be generated less frequently.

Practically, three compile-time routines allow to take care of the management of the $WOST\%$ garbage collection information :

- (1) *PROTECT* will check whether a new value to be stored on $WOST\%$ has to be protected through $GCTAB$; formulas given in II remain valid to make up that decision. When the value has to be protected through $GCTAB$, $gc1$ in $BOST$ and gcp in $BLOCKTAB$ have to be updated.
- (2) *DELETE* will cancel the protection of a given value deleted from $WOST\%$ by updating gcp in $BLOCKTAB$.
- (3) *UPDATE* will generate $gcw\%$ updating code if it appears that on $BLOCKTAB$ $gci \neq gcp$. This action is only to be taken when an ICI risking to call the garbage collector is to be generated. In $BLOCKTAB$, gci is updated accordingly.

APPENDIX 2 : SUMMARY OF THE SYNTAX

- 1. LBLOCK → lblockV BLOCKBODY
- 2.1 IDEDEC → idedecV FDECLARER iden = ACPAR
OPDEC → opdecV FDECLARER oper = ACPAR
- 2.2 LOCVARDEC → locvardecV ADECLARER variable |
locvardecV ADECLARER variable := SOURCE
- 2.3 HEAPVARDEC → heapvardecV ADECLARER variable
HEAPVARDEC → heapvardecV ADECLARER variable := SOURCE
- 3. LOCGEN → locV ADECLARER
HEAPGEN → heapV ADECLARER
- 4. LABELDEC → labeldecV label :
GOTO → gotoV label
- 5. CALL → callV PRIMCALL (ACPAR1, ACPAR2, ..., ACPARn)
FORMULA → dformulaV operator OPERAND1 OPERAND2 |
mformulaV operator OPERAND
ROUTDEN → routdenV (FORPAR1, FORPAR2,... FORPARn) : ROUTBODY
FORPARI → FDECLARER fideni
- 6. DEPROC → deprocV DEPROCCOERCEND
PROC → procV ROUTBODY
- 7. JPROC → jprocV label
- 8. CALLMODIND → callmodindV modind
MODEDEC → modedecV modind = ADECLARER
- 9. TRANSFORMAT → transformatV FORMATCOERCEND
FORMAT → formatV DYNREP
- 11.1 SELECTION → selectionV selector of SECONDARYSEL
- 11.2 Deref → derefV DEREFCOERCEND
- 11.3 SLICE → sliceV PRIMSLICE INDEXERS
INDEXERS → [INDEXER1 , ..., INDEXERn]
INDEXERi → TRIMMER |
INDEX
- 11.4 UNITED → unitingV UNCOERCEND
- 11.5 ROWING → rowingV ROWCOERCEND

- 12.1 ASSIGNATION \rightarrow assignationV DESTINATION := SOURCE
- 12.2 IDREL \rightarrow idrelV TERTL {:=: | :#:}₁¹ TERTR
- 12.3 CONFREL \rightarrow confrelV TERTL {:=: | ::}₁¹ TERTR
13. STDCALL \rightarrow stdcallV (ACPAR1, ACPAR2, ..., ACAPRn)
- 14.2 SERIAL \rightarrow LBLOCK | NONBLOCK
- LBLOCK \rightarrow lblockV BLOCKBODY
- BLOCKBODY \rightarrow NONBLOCK
- NONBLOCK \rightarrow SNONBLOCK | BALNONBLOCK
- SNONBLOCK \rightarrow PRELUDE lastV LABUNIT
- PRELUDE \rightarrow {{DECLA ; | UNITV ; }₀ ^{∞} DECLA ; }₀¹ {LABUNITVS}₀¹
- BALNONBLOCK \rightarrow {PRELUDE lastV LABUNIT . LABELDEC}₀¹
 {LABUNITVS lastV LABUNIT . LABELDEC}₀ ^{∞}
 LABUNITVS lastV LABUNIT
- LABUNITV \rightarrow {LABELDEC}₀ ^{∞} UNITV
- LABUNIT \rightarrow {LABELDEC}₀ ^{∞} UNIT
- LABUNITVS \rightarrow {LABUNITV ; }₀ ^{∞}
- DECLA \rightarrow IDEDEC | LOCVARDEC | HEAPVARDEC | OPDEC | MODEDEC
- 14.3 CONDCL \rightarrow ifV SERIALB CHOICECL fi
- CHOICECL \rightarrow then1V SERIAL |
 thefV SERIALB CHOICECL |
 then2V SERIAL1 elseV SERIAL2 |
 then3V SERIAL elsifV SERIALB CHOICECL
- 14.4 CASECL \rightarrow caseV CASECHOICE inV UNIT1, ... UNITn {outV SERIAL}₀¹ esac
- CASECHOICE \rightarrow UNITC | CASECONF
- 14.5 CASECONF \rightarrow caseconfV mode1V TERTL1, ... modenV TERTLn {:: | ::=}₁¹ TERTR
- 15.1 COLLVOID \rightarrow collvoidV (UNITV1, ..., UNITVn)
- 15.2 COLLROW \rightarrow collrowV (UNITD1, ..., UNITDn)
- 15.3 COLLSTR \rightarrow collstrV (UNITD1, ..., UNITDn)
- 16.3 FORCL \rightarrow forV {fromV UNITF}₀¹ {byV UNITB}₀¹
 {toV UNITT}₀¹ {foriden}₀¹ {whileV SERIALW}₀¹
 doV UNITD
- 16.4 TRCALL \rightarrow trcallV TRPRIM (UNIT1, ..., UNITn)

APPENDIX 3 : SUMMARY OF TOPST PROPERTIES

The table below shows how the fields *flextop* and Δmem of a TOPST element are initialized when a new TOPST element is set up by the activation of $p(\pi\alpha)$ or *NEWACTION* (*action*). Other TOPST fields are initialized to 0. The letter T means that the corresponding field in the new element is copied from the old top one.

$\pi(\alpha)$ /action	<i>flextop</i>	Δmem
BLOCKBODY	T	0
FDECLARER	(<i>stat</i> 0)	0
ACPARI	(<i>stat</i> 1)	0
OPERANDi	"	0
ADECLARER	(<i>stat</i> 0)	0
SOURCE	(<i>stat</i> 1)	0
PRIMCALL	(<i>stat</i> 0)	0
FORPARI	"	0
ROUTBODY	(<i>dyn bn</i>)	0
DEPROCCOERCEND	(<i>stat</i> 0)	0
FORMATCOERCEND	"	0
DYNREP	"	0
SECONDARYSEL	T	T
DEREFCOERCEND	(<i>stat</i> 0)	0
PRIMSLICE	"	T
TRIMMER	"	0
INDEX	"	0
UNCOERCEND	(<i>stat</i> 1)	T+ $\Delta union$
ROWCOERCEND	(<i>stat</i> 1)	T+ Δrow
DESTINATION	T	T
SOURCE	(<i>stat</i> 1)	T
TERTL	(<i>stat</i> 0)	0
TERTR	"	0
UNITV{i}	"	0
SERIALB	"	0
UNITC	"	0
TERTLi	"	0
<i>collrow</i> ∇	T	T
UNITDi	(<i>stat</i> 1)	0
<i>collstr</i> ∇	T	T
UNITF	(<i>stat</i> 0)	0
UNITB	"	0
UNITT	"	0
UNITD	"	0
SERIALW	"	0
TRPRIM	"	0
UNITi	"	0

APPENDIX 4 : SUMMARY OF THE NOTATIONS

1. MEM%

RANST% (*ranstpm%*)

HEAP% (*heappm%*)

DISPLAY%

BLOCK% | SBLOCK%
DBLOCK%

	Device	Static size	Current static pointers
S B L O C K %	H%	<i>h</i>	
	SIDST%	<i>sidsz</i>	<i>side</i>
	DMRWOST%	<i>dmrsz</i>	<i>dmrc</i>
	GCWOST%	<i>gcsz</i>	<i>gcc</i>
	SWOST%	<i>swostsz</i>	<i>swoste</i>
D B L O C K %	DIDST+LGST%		
	DWOST%		

2. H%

<i>stch%</i>	}	lblocks
<i>doh%</i>		
<i>wp%</i>		
<i>bn%</i>		
<i>gcid% gcidp%</i>	}	pblocks
<i>gebodyflag%</i>		
<i>gcw% gchp%</i>		
<i>gcsz%</i>	}	pblocks only
<i>result% swostp%</i>		
<i>gcp%</i>		
<i>dmrp%</i>		
<i>flex%</i>		
<i>prevflag%</i>		
<i>retjump%</i>		

3. DYNAMIC VALUE REPRESENTATION

Name : *pointer%*
scope%

Rowname : *pointer%*
scope%
descr%

Descriptor : *offset%*
states%
iflag%
do%
{i%
ui%
*di%}**

Union : *overhead%*
value%

Routine : *constabp%*
scope%

Format : *constabp%*
scope%

Tamrof : *offset%*
ndrep%
constabp%

4. BLOCKTAB (entry : *bnc*)

BLOCK%	Pseudo-BLOCK% _a	Routine BLOCK% _b
<i>sidsz</i>	$sidsz_a = sidsz_{b1}$	$sidsz_b = sidsz_{b1} + sidsz_{b2}$
<i>dmrsz</i>	$dmrsz_a$	$dmrsz_b$
<i>gcsz</i>	$gcsz_a$	$gcsz_b$
<i>swostsz</i>	$swostsz_a$	$swostsz_b$
<i>gcid</i>	$gcid_a$	$gcid_b$ { <i>gcbodyflag</i> }
<i>bn</i>	bn_a	bn_b

5. ACCESS (cadd)

<u>Fundamental</u>	<u>Accessory</u>
<u>(constant v)</u>	<u>(intet v)</u> <u>(boolet v)</u> <u>(bitset v)</u> <u>(charset v)</u>
<u>(directtab a)</u>	<u>(routet a)</u> <u>(formatet a)</u> <u>(tamrofet a)</u>
<u>(diriden bnc.side)</u> <u>(variden bnc.side)</u> <u>(indiden bnc.side)</u> <u>(dirwost bnc.swoste)</u> <u>(dirwost' bnc.swoste)</u> <u>(indwost bnc.swoste)</u> <u>(nihil 0)</u>	
	<u>(display bn)</u> <u>(dirabs a)</u> <u>(dirgcw bnc.gcc)</u> <u>(dirdmrw bnc.dmr)</u> <u>(varabs a)</u> <u>(varwost bnc.swoste)</u> <u>(i2iden bnc.side)</u> <u>(i2wost bnc.swoste)</u> <u>(label bnc.labnb)</u>

6. SYMBTAB

Identifiers (IDENTAB)

Static property	Fields	Form
<i>mode</i>		
<i>cadd</i>	<i>class</i> <i>add hadd</i> <i>tadd</i>	<i>(constant v)</i> <i>(directtab a)</i> <i>(diriden bnc.side)</i> <i>(variden bnc.side)</i> <i>(label bnc.labnb)</i>
<i>scope</i>	<i>inse</i> <i>outse</i>	
<i>flagdecl</i> <i>flagused</i>		

Mode indications (INDTAB)

<i>mode</i>		
<i>cadd</i>		<i>(label bnc.lo)</i>

7. BOST

Static property	Fields	Form
<i>mode</i>		<i>dectabp</i>
<i>cadd</i>	<i>class</i> <i>add hadd</i> <i>tadd</i>	
<i>smr</i>	<i>hadd</i> <i>tadd</i>	<i>bnc.swostc</i>
<i>dmr</i>		<i>(stat bnc.swostc)</i> <i>(dyn bnc.dmrce)</i> <i>nil</i>
<i>ge</i>		<i>bnc.gcc</i> <i>nil</i>
<i>or</i>	<i>kindo</i>	<i>iden</i> <i>var</i> <i>gen</i> <i>nil</i>
	<i>bno</i> <i>derefo</i> <i>geno</i> <i>{flexo}</i> <i>{diago}</i>	
<i>scope</i>	<i>inse</i> <i>outse</i>	
<i>{flexbot}</i>		
<i>obprogp</i>		

8. TOPST

Action	Fields	Form
<i>flextop</i>	<i>class</i> <i>spec</i>	(<i>stat 0</i>) (<i>stat 1</i>) (<i>dyn bn</i>)
<i>Δmem</i> <i>countbal</i> <i>countelem</i> <i>flagnextbal</i>		

9. CONSTAB

Routines

Non-standard	Standard	Jump
<i>lo</i> <i>bnsc</i> <i>sidesz_b</i> <i>dmrsz_b</i> <i>gcsz_b</i> <i>swostsz_b</i> <i>gcid_b</i> <i>flagstand {0}</i> <i>flagjump {0}</i>	{specific} <i>flagstand {1}</i>	<i>lo</i> <i>bnsc</i> <i>flagjump {1}</i>

Formats

lo
ndrep
bnsc
formstringp

APPENDIX 5 : LIST OF INTERMEDIATE CODE INSTRUCTIONS

This appendix is a complete list of the ICI's. With each of them, a number is given between brackets ; this number is the page number where the definition of the ICI is found.

1	STWOST1	MODE CADD CADD	(105)
2	STWOST2	MODE CADD CADD	(105)
3	STWOST3	MODE CADD CADD	(105)
4	STWORD	CADD CADD	(97)
5	STADD	CADD CADD	(104)
6	STGCWOST	MODE CADD CADDGC	(94)
7	STDMRWOST	CADD CADDMR	(106)
8	STACPAR	MODE CADD CADD	(109)
9	STDYNWOST1	MODE CADD	(197)
10	STDYNWOST2	MODE CADD	(197)
11	STDYNWOST3	MODE CADD	(197)
12	STSTATWOST	MODE CADD CADD	(105)
13	STGCNIL	CADDGC	(97)
14	PLUS	CADD CADD	(178)
15	STNDESCRWOST	MODE CADD	(184)

16	STGCELEM	BNC GCCS BNC0 GCCO	(254)
17	STOVERHUNION	MODE CADD	(195)
18	STPLUS	CADD1 CADD2 CADD0	(177)
19	STWOSTINCR	MODE CADD5 INCR CADD0	(254)
20	STNAMEINCR	CADD5 INCR CADD0	(179)
21	MINUS	CADD5 CADD0	(254)
22	STWP	BNC CADD	(241)
23	STOREREG	MODE CADD	(224)
24	LOADREG	MODE CADD	(224)
25	INCRRTWOSTPM	CADDINCR	(199)
26	HOLE		(222)
27	JUMP	LABNB	(139)
28	LABDEF	LABNB	(139)
29	JUMPNO	LABNB CADD	(214)
30	JUMPYES	LABNB CADD	(254)
31	GOTO	BNC BNCID LABNBID SWOSTC	(122)
32	LABFORMAT	CONSTABP	(254)
33	UPDCONSTAB	MODE CONSTABP	(142)
34	CHECKSTAND	LABNB CADD	(136)

35	LABID	LABNB	(121)
36	CHECKLAB	LABNB CADD	(148)
37	SWITCHCASE	LABNB CADD1 CADD2	(230)
38	OBPROG	OBPROGP	(254)
39	BALTAR	BALTABP	(254)
40	INPROG		(254)
41	OUTPROG		(254)
42	INBLOCK	BNC	(102)
43	CALLMIND	LRETURN BNCRES SWOSTCRES LBODY	(162)
44	OUTMIND	BNCBODY N CADD1 ***** CADDN	(163)
45	INMIND	BNCBODY	(163)
46	CALLDYNREP	LRETURN BNCRES SWOSTCRES CADDFORMAT	(167)
47	INDYNREP	BNCBODY	(170)
48	OUTDYNREP	BNCBODY N CADD1 ***** CADDN FORMATSTRINGP	(171)
49	INITDYNREP	CADDFORMAT CADDREP	(169)
50	INACPAR	BNCA FLEX CADDR0UT CADDRS GCCRES DMRCRES	(130)
51	CHECKDYNREP	LABNB CADDFORMAT	(169)

52	CALL	LRETURN CADDROUT BNCA	(132)
53	RETURN	MODERES CADDRS BNBODY	(140)
54	STANDCALL	NPAR DMRREC CADDROUT CADD1 ***** CADDN CADDRS DMRRES GCRES FLEX	(217)
55	STANDCALL1	LRETURN N BNCA CADDROUT CADD1 ***** CADDN	(137)
56	DEPROC	LRETURN CADDROUT CADDRS GCCRES DMRCRES FLEX	(146)
57	STANDDEPROC	LRETURN CADDROUT CADDRS GCCRES DMRCRES FLEX	(150)
58	CALLLAB	BNC CADDROUT	(157)
59	STDCALLINOUT	CADDROUT N MODE1 CADD1 ***** MODEN CADDN	(253)
60	CHECKFORMAL	MODE CADD N CADD1 ***** CADDN	(110)
61	CHECKFLEX	CADD	(186)

62	CHECKFLEXR	BNCROUT CADD	(186)
63	TRIMMER	NXDIM CADD CADDL CADDU CADDL' CADDOFF CADDT	(187)
64	INDEX	NXDIM CADD CADDI CADDOFF	(188)
65	STINTERSTFL	CADDFLAG CADDDESCR	(254)
66	STFILLSTRIDE	MODE CADDDESCR	(188)
67	STNAME	CADDPOINTER CADDSCOPE CADD0	(191)
68	FILLSTATEONE	CADDDESCR	(254)
69	STSCOPE	CADD CADD0	(254)
70	ROWINGSCADES	MODE0 CADD CADD0	(202)
71	ROWINGVAR	MODE0 CADD CADD0	(203)
72	ROWINGSCAL1	MODE0 CADD CADD0	(204)
73	ROWINGSCAL2	MODE0 CADD CADD0 DMRCS	(203)
74	ROWINGROW	MODES MODE0 CADD CADD0	(200)
75	ROWINGREFSCA	MODE0 CADD CADD0	(206)

76	ROWINGREFROW	MODES MODEO CADD5 CADD0	(205)
77	ROWINGEMPTY	MODE CADD	(175)
78	CHECKBOUNDS	CADDL CADDU CADDT	(240)
79	STBOUNDS	CADDL CADDU CADDT	(239)
80	STOVERHDESCR	MODEO STATES CADD CADD0	(239)
81	STFIRSTCOLLR	MODES MODEO CADD5 CADD0	(240)
82	STNEXTCOLLR	MODES MODEO CADD5 CADD0	(241)
83	STLITERALROW	MODE CADD5 CADD0	(202)
84	ROWS	MODEPAR2 CADDR0UT CADDPAR1 CADDPAR2 CADDRES	(254)
85	ASSIGN	MODE CADD5 CADD	(209)
86	ASSIGNSCOPE	MODE CADD5 CADD	(209)
87	LOCVARGEN	MODE CADD N CADD1 CADDN	(112)

88	HEAPVARGEN	MODE CADD N CADD1 ***** CADDN	(115)
89	FORTO	LABNB CADDFOR1 CADDBY CADDTO	(250)
90	CHECKSCBLOCK	MODE CADD BN	(104)
91	IDREL=	MODESL CADDSL CADDSR CADD0	(211)
92	IDREL =	MODESL CADDSL CADDSR CADD0	(211)
93	LOGGEN	MODE CADD CADDGC N CADD1 ***** CADDN	(118)
94	HEAPGEN	MODE CADD CADDGC N CADD1 ***** CADDN	(119)
95	CONFTO	MODEL MODER CADDR CADD0	(213)
96	CONFTOBEC	MODEL MODER CADD0 CADDR GCR DMRR CADD*R GC*R DMR*R	(213)

97	WIDEN	MODES MODEO CADD CADD	(248)
98	STNIL	CADD	(174)
99	STSKIP	MODE CADD	(173)
100	CONJWOST	CADD	(215)
101	CONJ	CADD CADD	(215)
102	NOOPTIMIZE		(94)
103	NEWCARD	CARDNB	(254)
104	PRID	IDEN	(254)
105	PRNUMB	NUMB	(254)
106	CHECKNIL	CADD	(179)
107	CHECKOVERLAP	MODE CADD CADD LABNB	(209)
108	STSTATACPAR	MODE CADD CADD	(109)
109	DEPROC1	LRETURN CADDR CADDR GCCRES DMRCRES FLEX	(152)
110	INBODY	BNCBODY	(154)

APPENDIX 6 : AN EXAMPLE OF COMPILATION

This example has been chosen very simple, it is only intended to give a flavour on how the successive stages of the compilation look like. (More examples can be found in [21]). The following comments on the output from the computer are useful :

- (1) The source program text is first printed, its lines are numbered. The numbering, referred to as card number, is used in the next outputs as a reference to the source program.
- (2) `DECTAB$` and
- (3) `IDENTAB$` are self explaining. Note that the numbering in the first column is used for cross referencing.
- (4) `BLOCKTAB$` must be explained :
 - line 25100 : corresponds to the block 'program'.
 - line 25103 : corresponds to 'particular program' with
 - `bn=1` lexicographical depth number
 - `sidsz=2` resulting from the way `Q(5)` is translated : the copy of the value of the primary `Q` is forced on `WOST%`, moreover, space is foreseen for the result of the call.
 - `head` and `tail` are `IDENTAB$` pointers, keeping track of the declarations of the block, chained together ; they play the role of *gold*.
 - line 25106 : corresponds to the routine possessed by `P` with
 - `bn=1`, i.e. the scope 0 of the routine +1.
 - `sidsz=2` corresponding to parameter `X` and variable `A` (blocks are merged).
 - `swostsz=1`, foreseen for the result of `10+X`.
 - line 25109 : corresponds to the pseudo-block of the actual parameter of the call `Q(5)`, with
 - `bn=2`, lexicographical depth number.
 - `sidsz=1` for the integer parameter.
- (5) Follows the linear prefixed form of the program i.e. `SOPROG` for the IC generation. This form is self-explaining. Note however that
 - prefix markers start with '\$',
 - the lines `**** CARD` refer to the source program (1)
 - the numbers 235, 236,... respectively refer to `IDENTAB$` entries 30584, 30588,...
 - the number 232 refers to the operator `op(int,int)int +` in the initialized part of `INDTAB$` not printed here.
 - coercions are kept in a separate table `COERCTAB` ; connections with `SOPROG`

are obtained through the specification field of the prefix markers of the coerecends ; here : \$ID and \$DEN.

- (6) The intermediate code (*OBPROG*) should be easily understood in the light of PART II of this book. Only some details may differ from what has been described. Note moreover that *CONSTAB* referred to in some ICI's is not printed here.
- (7) The machine code in its relocatable form is then printed ; in this code, we find :
- (a) instruction lines consisting of :
- the opcode using mnemonics. Note that A, S, G and F are registers and that SUBC means a subroutine call. The opcode may be preceded by one letter U, Y or N to mean a conditional execution of the instruction.
 - an optional star the presence of which means (hardware) literal addressing.
 - STAT, MPQ, MA and MS which are addressing types :
 - STAT means normal addressing.
 - MPQ means (hardware) display addressing.
 - MA (MS) means indexed addressing using A (S) register.
 - the address field consisting of a pair of integers, the first one being only significant in case of display addressing. This field may be followed by P, Z or E, which causes the setting up of conditions, subsequently to the execution of the instruction.
 - finally the field symbolic which is a symbolic representation of run-time routines or static working cells. In case it is 'LABTABPI' however, it has a special meaning, indicating to the loader that address conversion using *LABTAB* is involved.
- (b) loader commands :
- LABDEF for label definition.
 - LABROUT and OFFSECT issued from the IC *updconstab* with mode parameters proc and string respectively.
- (c) lines with '****' :
- They correspond to references to the ICI's currently translated.
- (8) The loader automatically prints the starting addresses of :
- the object program *OBPROG*,
 - *RANST* and
 - *HEAP*
- (9) The actual result of the program execution is finally printed.

1. SOURCE PROGRAM

```

1      (
2          PROC P=(INT X)INT :(INT A:=10+X;
3              ----      ---      ---      ---
4                      A);
5          PR 51PR  PR 62PR  PR 63PR  PR 64PR
6          --  --  --  --  --  --  --  --
7          PROC (INT )INT Q:=P;
8          ----  ---  ----
9          PRINT(("RESULT=",Q(5)))
10         )

```

2. DECTAB

```

15500  PROC
        MODERES      INT
        NMBPAR        1
        PARAM1        INT
15503  PROC
        MODERES      INT
        NMBPAR        1
        PARAM1        INT
15506  REF            15503

```

3. IDENTAB

	<i>cadd</i>	<i>used</i>	<i>mode</i>	<i>scope</i>	<i>chain</i>	<i>length</i>	<i>alpha</i>
30584	ROUTCT 0 1342	1	15500	0 0	0	1	P
30588	DIRIDEN 2 0	1	INT	0 0	30592	1	X
30592	VARIDEN 2 1	1	REF INT	1 1	0	1	A
30596	VARIDEN 1 0	1	15506	1 1	0	1	Q

4. BLOCKTAB

	<i>swostsz</i>	<i>sidez</i>	<i>bn</i>	<i>head</i>	<i>gsz</i>	<i>tail</i>	<i>dmrsz</i>
25100	0	0	0	0	0	0	0
25103	3	2	1	30596	0	30592	0
25106	1	2	1	30588	0	30592	0
25109	0	1	2	0	0	0	0

5. SOPROG

0	****CARD	1	53	\$ID	0
1	****CARD	2	54	ID	152
2	(CLO	2	55	(AP	0
3	<RAN	0	56	(COLL	1
4	\$CONSD	0	57	COLLMOD	14555
5	DECLAR	15500	58	\$DEN	5
6	DEFID	235	59	STRINGT	0
7	(R	0	60	,	0
8	SCOPE	0	61	\$CALL	7
9	DECLAR	15500	62	\$ID	3
10	(F	0	63	ID	238
11	DECLAR	14570	64	(AP	0
12	DEFID	236	65	\$DEN	0
13)F	0	66	SINTCT	5
14	DECLAR	14570	67)PA	0
15	:	0	68)LLOC	1
16	(CLO	1	69)PA	0
17	<RAN	0	70	****CARD	7
18	\$LVARDE	0	71	>NAR	0
19	DECLAR	14570	72)OLC	2
20	DEFID	237			
21	***	0			
22	\$FORMUL	0			
23	DYADOPE	232			
24	\$DEN	0			
25	SINTCT	10			
26	SEP	0			
27	\$ID	0			
28	ID	236			
29	;	0			
30	\$LASTUN	0			
31	****CARD	3			
32	\$ID	1	0	END	0
33	ID	237	1	DEREF	0
34	>NAR	0	2	END	14587
35)OLC	1	3	DEREF	0
36)R	0	4	END	15506
37	;	0	5	[]OUTTYPE	0
38	\$LVARDE	0	6	END	14660
39	DECLAR	15503	7	[]OUTTYPE	0
40	****CARD	4			14570
41	PRAGNUMB	62			
42	PRAGNUMB	63			
43	****CARD	5			
44	PRAGNUMB	64			
45	DEFID	238			
46	***	0			
47	\$ID	0			
48	ID	235			
49	;	0			
50	\$LASTUN	0			
51	\$CALL	0			
52	****CARD	6			

COERCTAB

6. OBPROGS

0	INPROG				
1	****CARD1				
2	****CARD2				
3	INBLOCK	BNC			1
4	JUMP	LABNB			3
5	C2:				
6	LOCVARGEN	MODE	INT		
		CADD	DIRIDEN	2	1
		N			0
9	STANDCALL	LRETURN			0
		NLONG			0
		N			2
		BNC			2
		DMRC	NIL*		
		CADDROUT	ROUTCT	0	1129
		CADD 1	INTCT	0	10
		CADD 2	DIRIDEN	2	0
		CADDRS	DIRWOST	2	0
		DMRCRES	NIL*		
		GCCRES			0
		FLEX	STAT		1
17	ASSIGN	MODES	INT		
		CADD5	DIRWOST	2	0
		CADD0	VARIDEN	2	1
20	****CARD3				
21	HOLE				
22	UPDCONSTAB	MODE			15500
		CADD	ROUTCT	0	1342
24	RETURN	MODERES	INT		
		CADDRS	DIRIDEN	2	1
		BNBODY			1
27	C3:				
28	****CARD4				
29	PRAGMAT	62			
30	PRAGMAT	63			
31	****CARD5				
32	PRAGMAT	64			
33	LOCVARGEN	MODE			15503

		CADD N	DIRIDEN	1	0 0
36	ASSIGN	MODES CADD CADD	ROUTCT VARIDEN	0 1	15500 1342 0
39	****CARD6				
40	UPDCONSTAB	MODE CADD	DIRCTTAB	0	14660 1345
42	STWOST3	MODE CADD CADD	DIRIDEN DIRWOST	1 1	15503 0 0
45	INACPAR	BNCACPA FLEX CADDROUT CADDRS GCCRSE DMRCRES	STAT DIRIDEN DIRWOST	1 1	3 0 0 2 0 0
49	STACPAR	MODE CADD CADD	INTCT DIRIDEN	INT 0 3	5 0
52	CHECKSTAND	LABNR CADD	DIRWOST	1	4 0
54	CALL	LRETURN CADDROUT BNCACPA	DIRWOST	1	5 0 3
57	C4:				
58	STANDCALL1	LRETURN N BNC CADDROUT CADD 1	DIRWOST DIRIDEN	1 3	5 1 3 0 0
62	C5:				
63	****CARD7				
64	STDCALLINOUT	LRETURN N BNC DMRC CADDROUT MODE 1 CADD 1 MODE 2 CADD 2 CADDRS DMRCRES GCCRES FLEX	NIL* ROUTCT DIRCTTAB DIRWOST NIHIL NIL* STAT	0 0 0 1 0 0 0	0 2 1 565 14660 1345 2 0 0 NILGC 0
75	HOLE				
76	STADD	CADD	DDISPLAY	0	1

		CADD0	DIRABS	0	20
79	STWORD	CADDS	INTCT	0	0
		CADD0	DIRABS	0	17
82	STWORD	CADDS	VARABS	0	4095
		CADD0	DDISPLAY	0	1
85	***CARD8				
86	L1:				
87	OUTPROG				

7. MACHINE CODE

	<i>opcode</i>	<i>lit</i>	<i>addtype</i>	<i>addr</i>	<i>addr</i>	<i>symb</i>
0	LDA	*	STAT	0	11	
1	LDS	*	STAT	0	0	
2	LDG	*	STAT	0	11	
3	STG		STAT	0	0	INCRGC9
4	LDG	*	STAT	0	0	
5	SUBC	*	STAT	0	0	INPROG9
	****				1	
	****				2	
	****				3	
6	LDG	*	STAT	0	2	
7	STG		STAT	0	0	CARDNB
8	LDA	*	STAT	0	1	
9	LDS	*	STAT	0	0	GCINFOTAB
10	LDG	*	STAT	0	16	
11	SUBC	*	STAT	0	0	INBLOCK19
12	LDA	*	STAT	0	13	
13	LDS	*	STAT	0	0	
14	LDG	*	STAT	0	2	
15	SUBC	*	STAT	0	0	INBLOCK29
	****				4	
	****				5	
16	GOTO	*	STAT	0	29	LABTABPI
	****				6	
	****				9	
	LABDEF				2	
17	LDS	*	STAT	0	10	
18	ADS		MPQ	1	11	
	****				17	
	****				20	
	****				21	
	****				22	
19	STS		MPQ	1	12	
	LABROUT				17	LABTABPI
	****				24	
20	LDG		MPQ	1	4	
21	STG		STAT	0	0	RET.JUMP
22	LDS		MPQ	1	5	
23	LDA	*	MPQ	1	12	
24	LDG	*	STAT	0	1	
25	SUBC	*	STAT	0	0	UPDATEDISP
26	LDG		MA	0	0	
27	STG		MS	0	0	
28	GOTO		STAT	0	0	RET.JUMP

	****				27	
	****				28	
	****				29	
	****				30	
	****				31	
	****				32	
	****				33	
	****				36	
	LABDEF				3	
29	LDG	* STAT	0	1342		CONSTABPI
30	STG	MPQ	1	11		
31	LDG	STAT	0	0		DISPLAYPI
	****				39	
	****				40	
32	STG	MPQ	1	12		
	OFFSECT			1345		
	****				42	
33	LDF	MPQ	1	11		
	****				45	
34	STF	MPQ	1	13		
35	LDG	* STAT	0	6		
36	STG	STAT	0	0		CARDNB
37	LDG	* STAT	0	1358		CONSTABPI
38	LDS	MPQ	1	11		
39	SUBC	* STAT	0	0		INACPARDYN
	****				49	
40	LDS	* STAT	0	5		
	****				52	
41	STS	MPQ	2	11		
42	LDA	MPQ	1	13	Z	
43	Y LDS	* STAT	0	14		
44	Y GOTO	* STAT	0	0		ALARM
45	U LDA	MA	0	0	P	
	****				54	
46	N GOTO	* STAT	0	53		LABTABPI
47	LDA	* STAT	0	57		LABTABPI
48	STA	STAT	0	0		RET.JUMP
49	LDA	* STAT	0	1		
50	LDS	* MPQ	1	13		
51	LDG	* STAT	0	2		
52	GOTO	* STAT	0	0		CALLDYN
	****				57	
	****				58	
	LABDEF				4	
53	LDG	* STAT	0	6		
54	STG	STAT	0	0		CARDNB
55	LDS	* STAT	0	21		
56	GOTO	* STAT	0	0		ALARM
	****				62	
	****				63	
	****				64	
	LABDEF				5	
57	SUBC	* STAT	0	0		SAVETIME
58	LDS	* STAT	0	1345		CONSTABPI
59	LDA	MS	0	2	Z	
60	Y GOTO	* STAT	0	70		LABTABPI
61	LDA	* STAT	0	0		VALSTPI
62	LDG	* STAT	0	1		
63	STG	STAT	0	2		VALSTPI
64	LDG	* STAT	0	1000		
65	SUBC	* STAT	0	0		INITLOOPR9
	LABDEF				6	
66	SUBC	* STAT	0	0		PRINTCHARS

67	LDB	*	STAT	0	66	LABTABPI
68	LDA	*	STAT	0	0	VALSTPI
69	SUBC	*	STAT	0	0	FINALLOOPI9
	LABDEF				7	
70	LDS	*	MPQ	1	15	
71	SUBC	*	STAT	0	0	PRINTINTS
72	SUBC	*	STAT	0	0	RESTTIME
	****				75	
	****				76	
73	LDA		STAT	0	1	DISPLAYPI
74	SBA	*	STAT	0	256	
	****				79	
75	STA		STAT	0	0	RTWOSTPM
76	LDS	*	STAT	0	0	
	****				82	
77	STS		STAT	0	0	RTBNA
78	LDS	*	STAT	0	0	NIL
	****				85	
	****				86	
79	STS		STAT	0	1	DISPLAYPI
	****				87	
	LABDEF				1	
80	GOTO	*	STAT	0	0	FINAL9

8. LOADER INDICATIONS

```

OBPROGPI 11368
STACKPI  11449
HEAPPI   29700

```

9. PROGRAM RESULT

```

RESULT=      15

```