A Labelled Transition System for π_ϵ-Calculus

Franck van Breugel

Università di Pisa, Department of Computer Science
Corso Italia 40, 56125 Pisa, Italy

Abstract. A labelled transition system is presented for Milner's π_ϵ-calculus. This system is related to the reduction system for the calculus presented by Bellin and Scott. Also a reduction system and a labelled transition system for π_ϵ1-calculus are given and their correspondence is studied. This calculus is a subcalculus of π_ϵ-calculus in the way Sangiorgi's π1-calculus is a subcalculus of ordinary π-calculus.

Introduction

In the early nineties, Abramsky [Abr94] presented a translation from proofs in linear logic into π-calculus, and outlined the results relating the computational behaviour of the proofs under cut-elimination to that of the processes under reductions. When Milner heard of Abramsky's result, he worked out his own translation. This led to the development of a synchronous version of π-calculus [Mil93], which we call π_ϵ-*calculus*[1]. In [BS94], Bellin and Scott analysed Abramsky's translation in detail for Milner's π_ϵ-calculus.

In π_ϵ-calculus we encounter enabling, extended scope extrusion, and self communication. These three features are not present in ordinary π-calculus. We discuss them in the following three paragraphs.

In π-calculus, the process $\alpha.P$ specifies that the action α has to precede all actions in P. For the π_ϵ-calculus process $\alpha\,P$ this condition has been weakened as follows. The action α only has to precede those actions in P which it *enables*, i.e. those actions a free name of which is bound by α. For example, in the process $w(x)\,\bar{y}z\,P$, where $x \neq y$, z, the action $w(x)$ does not enable $\bar{y}z$. As a consequence, the action $\bar{y}z$ may precede $w(x)$. Hence, if we put the process $w(x)\,\bar{y}z\,P$ in parallel with $y(z)\,Q$, then a communication at y can occur resulting in the process $w(x)\,P$ in parallel with Q. This is modelled by the reduction

$$w(x)\,\bar{y}z\,P \mid y(z)\,Q \rightarrow w(x)\,P \mid Q. \tag{1}$$

Like in π-calculus, in π_ϵ-calculus we encounter scope extrusions. For example, if $x \neq y$ then

$$(\nu x)\,\bar{y}x\,P \mid y(x)\,Q \rightarrow (\nu x)\,(P \mid Q). \tag{2}$$

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[1] Since we do not want to contrast the calculus with *asynchronous* π-calculus [HT91, Bou92] and enablement is one of its key features, we call it π_ϵ-calculus.

Usually, only scopes of the form (νx) are extruded. In π_ϵ-calculus also *extended scopes* like $(\nu w)\, w(x)\, x(y)$—a formal definition of these extended scopes is given in Definition 6—are extruded. For example,

$$(\nu w)\, w(x)\, x(y)\, \bar{z}y\, P \mid z(y)\, Q \to (\nu w)\, w(x)\, x(y)\, (P \mid Q) \tag{3}$$

provided that $z \neq w,\, x,\, y$, and $w,\, x$ do not occur free in Q.

In π_ϵ-calculus, a process can communicate with itself. In its simplest form, *self communication* amounts to

$$\bar{x}y\, x(z)\, P \to P[y\!/\!z]. \tag{4}$$

Self communication can also take place in extended scopes. For example, if $w \neq x,\, y,\, z$ then

$$w(x)\, (\nu y)\, y(z)\, \bar{w}z\, P \to (\nu y)\, y(z)\, (P[z\!/\!x]). \tag{5}$$

The process communicates with itself at w within the extended scope $(\nu y)\, y(z)$.

For π_ϵ-calculus, Bellin and Scott [BS94] presented a *reduction system* [Mil92]. We briefly review this system in Section 2. The rules defining this system are simple and natural. However, the system does not support reasoning in a purely structural way. In Section 3, we give a *labelled transition system* for the calculus following Plotkin's structural approach [Plo81]. The rules defining the labelled transition system are non-trivial. In Section 4, the correctness of this system is shown by proving the correspondence between the reduction system and the labelled transition system. Both the reduction system and the labelled transition system are useful (cf. [San92, page 26]) and once their relation has been established they support each other.

In [San96], Sangiorgi studied a subcalculus of π-calculus, called πI-calculus, which only uses *internal mobility*. In Section 5, we present a reduction system and a labelled transition system for π_ϵI-calculus, a subcalculus of π_ϵ-calculus with only internal mobility. Furthermore, we investigate the relation between the two systems.

Some related work is discussed in Section 6. In the final section, some conclusions are drawn. We assume that the reader is familiar with π-calculus and πI-calculus. For an introduction to π-calculus we refer the reader to Milner's tutorial [Mil91]. In Sangiorgi's [San96], πI-calculus is studied in great detail.

1   Basic π_ϵ-calculus

We assume an infinite set of *names*. We use x, y, x_1, y', ... to range over these names.

Definition 1. The set of *processes* is defined by

$$P ::= 0 \mid \pi\,P \mid P \mid Q$$

where the set of *particles* is given by

$$\pi ::= \bar{x}y \mid x(y) \mid (\nu x) \qquad\qquad \lrcorner$$

Only the constructs $\bar{x}y\,P$ and $x(y)\,P$ are not part of ordinary π-calculus. Just a small fragment of π_ϵ-calculus is presented here. We are confident that the results of the present paper can be extended straightforwardly if we add operators like summation and replication.

This calculus has two binders, the particles $x(y)$ and (νx). We define the *bound names* and *free names* of particles and processes in the usual way.

π	bn (π)	fn (π)
$\bar{x}y$	$\emptyset$	$\{x, y\}$
$x(y)$	$\{y\}$	$\{x\}$
(νx)	$\{x\}$	$\emptyset$

P	bn (P)	fn (P)
0	$\emptyset$	$\emptyset$
$\pi\,P$	bn (π) $\cup$ bn (P)	fn (π) $\cup$ (fn (P) \ bn (π))
$P \mid Q$	bn (P) $\cup$ bn (Q)	fn (P) $\cup$ fn (Q)

The *names* of particles and processes are given by n $(\pi) = $ bn $(\pi) \cup$ fn (π) and n $(P) = $ bn $(P) \cup$ fn (P).

2 Reduction system

The reduction system is defined in two steps. First, we identify several processes by introducing a structural congruence over processes. Second, we define the computation steps of processes in terms of a reduction relation. Our presentation is based on [Mil91, Section 2] and [BS94, Section 2].

Definition 2. The *structural congruence* $\equiv$ is defined as the smallest congruence relation over processes satisfying

1. if P and Q are alpha-convertible then $P \equiv Q$
2. $P \mid Q \equiv Q \mid P$
3. $(P \mid Q) \mid R \equiv P \mid (Q \mid R)$
4. $0 \mid P \equiv P$
5. if n $(\pi_1) \cap$ bn $(\pi_2) = \emptyset$ and n $(\pi_2) \cap$ bn $(\pi_1) = \emptyset$ then $\pi_1\,\pi_2\,P \equiv \pi_2\,\pi_1\,P$
6. if bn $(\pi) \cap$ fn $(Q) = \emptyset$ then $\pi\,(P \mid Q) \equiv (\pi\,P) \mid Q$ $\qquad \lrcorner$

For ordinary π-calculus 1., 2., 3., and 4., and 5. and 6. restricted to particles of the form (νx) are used (see [Mil91, page 7 and 8]). In [Eng96, page 81], Engelfriet considers the following variation of 5.

if $x \notin$ n (π) then $(\nu x).\pi.P \equiv \pi.(\nu x).P$

Definition 3. The *reduction relation* $\to$ is defined as the smallest relation over processes satisfying

1. $x(y)\,P \mid \bar{x}z\,Q \to P[z\!/\!y] \mid Q$

2. $\dfrac{P \to P'}{\pi\,P \to \pi\,P'}$

3. $\dfrac{P \to P'}{P \mid Q \to P' \mid Q}$

4. $\dfrac{P \equiv Q \qquad Q \to Q' \qquad Q' \equiv P'}{P \to P'}$

For ordinary π-calculus one only needs 1.,3., and 4. (see [Mil91, page 8]).

In [Bre97] we give proofs of the reductions presented in the introduction. We conclude this section with some properties of the structural congruence. These will be exploited when we link the reduction system and the labelled transition system.

Proposition 4. *If $P \equiv Q$ then $\mathrm{fn}\,(P) = \mathrm{fn}\,(Q)$.*

Proof. Induction on the proof of $P \equiv Q$. □

Proposition 5. *If $P \equiv Q$ then $P[x\!/\!y] \equiv Q[x\!/\!y]$.*

Proof. Induction on the proof of $P \equiv Q$. □

In Proposition 7 we show that 5. and 6. of Definition 2 also hold for scopes.

Definition 6. The set of *connected input sequences* is given by

$$\iota_x^y ::= x(y) \mid x(z)\,\iota_z^y$$

The set of *scopes* is defined by

$$\sigma_x ::= \iota_y^x \mid (\nu x) \mid (\nu y)\,\iota_y^x$$

A connected input sequence $\iota_{x_1}^{x_n}$ is of the form $x_1(x_2)\,x_2(x_3)\ldots x_{n-1}(x_n)$. These are related to Sangiorgi's dependency chains [San96, Definition 6.5]. In ordinary π-calculus one usually only considers scopes of the form (νx). The role of these extended scopes will be discussed in the next section. The bound and free names of scopes are defined straightforwardly.

ι	$\mathrm{bn}\,(\iota)$	$\mathrm{fn}\,(\iota)$
$x(y)$	$\{y\}$	$\{x\}$
$x(z)\,\iota_z^y$	$\{z\} \cup \mathrm{bn}\,(\iota_z^y)$	$\{x\}$

σ	$\mathrm{bn}\,(\sigma)$	$\mathrm{fn}\,(\sigma)$
ι_y^x	$\mathrm{bn}\,(\iota_y^x)$	$\mathrm{fn}\,(\iota_y^x)$
(νx)	$\{x\}$	$\emptyset$
$(\nu y)\,\iota_y^x$	$\{y\} \cup \mathrm{bn}\,(\iota_y^x)$	$\emptyset$

Proposition 7.

1. *If $\mathrm{n}\,(\sigma) \cap \mathrm{bn}\,(\pi) = \emptyset$ and $\mathrm{n}\,(\pi) \cap \mathrm{bn}\,(\sigma) = \emptyset$ then $\sigma\,\pi\,P \equiv \pi\,\sigma\,P$.*
2. *If $\mathrm{bn}\,(\sigma) \cap \mathrm{fn}\,(Q) = \emptyset$ then $\sigma\,(P \mid Q) \equiv (\sigma\,P) \mid Q$.*

Proof. Structural induction on σ. □

3 Labelled transition system

The labelled transition system presented in this section is new. Its presentation is based on [MPW92, page 46] and [ACS96, page 150]. The system uses the late scheme of name instantiation. It can be adapted straightforwardly to deal with the early scheme (cf. [MPW91, page 49 and 50]).

The labelled transition system not only describes the computation steps of processes but also their communications possibilities. This information is recorded by means of actions.

Definition 8. The set of *actions* is given by

$$\alpha ::= \bar{x}y \mid x(y) \mid \tau \mid \sigma_y \, \bar{x}y$$

where $x \notin \text{bn}\,(\sigma_y)$. ⌟

In ordinary π-calculus the action $(\nu y)\,\bar{x}y$ is usually written as $\bar{x}(y)$. The actions $\sigma_y \, \bar{x}y$, with $\sigma_y \neq (\nu y)$, one does not encounter in the usual labelled transition system. These extended scopes are used to model extended scope extrusions (cf. (3) in the introduction).

The bound and free names of actions are defined as follows.

α	$\text{bn}\,(\alpha)$	$\text{fn}\,(\alpha)$
$\bar{x}y$	$\emptyset$	$\{x,y\}$
$x(y)$	$\{y\}$	$\{x\}$
τ	$\emptyset$	$\emptyset$
$\sigma_y \, \bar{x}y$	$\text{bn}\,(\sigma_y)$	$\{x\}$

In the next definition the transition relation is presented. We have omitted the symmetric versions of the rules 9., 10., and 11.

Definition 9. The *transition relation* $\to$ is defined as the smallest labelled relation over processes satisfying

1. $\dfrac{P' \xrightarrow{\alpha} Q'}{P \xrightarrow{\alpha} Q}$ $\quad P$ and P', and Q and Q' are alpha-convertible

2. $\bar{x}y\, P \xrightarrow{\bar{x}y} P$

3. $x(y)\, P \xrightarrow{x(y)} P$

4. $\dfrac{P \xrightarrow{\alpha} P'}{\pi\, P \xrightarrow{\alpha} \pi\, P'}$ $\quad \text{n}\,(\alpha) \cap \text{bn}\,(\pi) = \emptyset$ and $\text{n}\,(\pi) \cap \text{bn}\,(\alpha) = \emptyset$

5. $\dfrac{P \xrightarrow{\bar{z}y} P'}{x(y)\, P \xrightarrow{x(y)\,\bar{z}y} P'}$ $\quad y \neq z$ $\qquad \dfrac{P \xrightarrow{\iota_y^w \, \bar{z}w} P'}{x(y)\, P \xrightarrow{x(y)\,\iota_y^w \, \bar{z}w} P'}$ $\quad y \neq z$

6.
$$\frac{P \xrightarrow{\bar{z}y} P'}{(\nu y)\, P \xrightarrow{(\nu y)\,\bar{z}y} P'} \quad y \neq z \qquad \frac{P \xrightarrow{\iota_y^w \,\bar{z}w} P'}{(\nu y)\, P \xrightarrow{(\nu y)\,\iota_y^w \,\bar{z}w} P'} \quad y \neq z$$

7.
$$\frac{P \xrightarrow{x(z)} P'}{\bar{x}y\, P \xrightarrow{\tau} P'[y/z]}$$

8.
$$\frac{P \xrightarrow{\bar{x}z} P'}{x(y)\, P \xrightarrow{\tau} P'[z/y]} \quad y \neq x, z \qquad \frac{P \xrightarrow{\sigma_z\,\bar{x}z} P'}{x(y)\, P \xrightarrow{\tau} \sigma_z\,(P'[z/y])} \quad y \notin \mathrm{n}\,(\sigma_z\,\bar{x}z)$$

9.
$$\frac{P \xrightarrow{\alpha} P'}{P \mid Q \xrightarrow{\alpha} P' \mid Q} \quad \mathrm{bn}\,(\alpha) \cap \mathrm{fn}\,(Q) = \emptyset$$

10.
$$\frac{P \xrightarrow{x(z)} P' \qquad Q \xrightarrow{\bar{x}y} Q'}{P \mid Q \xrightarrow{\tau} P'[y/z] \mid Q'}$$

11.
$$\frac{P \xrightarrow{x(y)} P' \qquad Q \xrightarrow{\sigma_y\,\bar{x}y} Q'}{P \mid Q \xrightarrow{\tau} \sigma_y\,(P' \mid Q')} \quad \mathrm{bn}\,(\sigma_y) \cap \mathrm{fn}\,(P) = \emptyset$$

Some remarks:

- The rules 1., 4. with π of the form (νx), the first part of 6., 9., 10., and 11. with σ_y of the form (νy) are as usual.
- The axioms 2. and 3. are as expected.
- The rule 4. models enabling and corresponds to Definition 2.5.
- The rules 5. and 6. describe *scope opening* (cf. [MPW92, page 48]). Like in ordinary π-calculus, the side condition $y \neq z$ prevents z from becoming bound (cf. Definition 8). The rule 11. handles *scope closing*. Note that the scope σ_y reappears in the conclusion. The side condition $\mathrm{bn}\,(\sigma_y) \cap \mathrm{fn}\,(P) = \emptyset$ prevents us from deriving the incorrect transition

$$x(y)\,\bar{z}y\,0 \mid (\nu z)\,z(y)\,\bar{x}y\,0 \xrightarrow{\tau} (\nu z)\,z(y)\,(\bar{z}y\,0 \mid 0).$$

This transition is incorrect since the free name z in $x(y)\,\bar{z}y\,0$ is only accidentally the same as the bound name z in $(\nu z)\,z(y)\,\bar{x}y\,0$. In [Bre97] the interplay between scope opening and scope closing is illustrated.

- The rules 7. and 8. describe self communication. Because of the side condition $y \neq x$, we cannot prove the obviously incorrect transition

$$x(y)\,\bar{y}z\,0 \xrightarrow{\tau} 0.$$

This side condition ensures that the free name x and the bound name y, which can be alpha-converted to x, are not identified. The side condition $y \neq z$ rules out the transition

$$x(y)\,\bar{x}y\,\bar{x}y\,0 \xrightarrow{\tau} \bar{x}w\,0.$$

This transition is incorrect because the bound name y, which can be alpha-converted to w, becomes free (cf. Proposition 10). The side condition $y \notin \text{n}\,(\sigma_z\,\bar{x}z)$ prevents us from proving the transition

$$x(y)\,(\nu y)\,y(z)\,\bar{x}z\,\bar{y}w\,0 \xrightarrow{\tau} (\nu y)\,y(z)\,\bar{z}w\,0.$$

This transition is incorrect since the name y in $\bar{y}w$ is bound by (νy) in the process $x(y)\,(\nu y)\,y(z)\,\bar{x}z\,\bar{y}w\,0$ whereas the corresponding z in $\bar{z}w$ is bound by $y(z)$ in $(\nu y)\,y(z)\,\bar{z}w\,0$.

– The following variation of 1. (cf. [San92, page 30]) suffices to prove the results of Section 4.

$$\frac{P' \xrightarrow{\alpha} Q}{P \xrightarrow{\alpha} Q} \quad P \text{ and } P' \text{ are alpha-convertible}$$

We have chosen for 1. since it is convenient for proving Proposition 11.

In [Bre97] we give proofs of the τ-transitions corresponding to the reductions presented in the introduction. Like in the previous section, we conclude with some properties which are used when we relate the two systems.

Proposition 10. *If* $P \xrightarrow{\alpha} P'$ *then* $\text{fn}\,(P') \subseteq \text{fn}\,(P) \cup \text{bn}\,(\alpha)$ *and* $\text{fn}\,(\alpha) \subseteq \text{fn}\,(P)$.

Proof. Induction on the proof of $P \xrightarrow{\alpha} P'$. $\qquad\qquad\square$

Proposition 11. *If there exists a proof of* $P \xrightarrow{\alpha} Q$ *not containing* $\text{bn}\,(\alpha')$ *and* $\alpha\,Q$ *and* $\alpha'\,Q'$ *are alpha-convertible[2] then* $P \xrightarrow{\alpha'} Q'$.

Proof. Induction on the depth of the proof of $P \xrightarrow{\alpha} Q$. $\qquad\qquad\square$

4 Correspondence between the systems

The reduction system of Section 2 and the labelled transition system of Section 3 are related in this section. More precisely, reductions and τ-transitions are linked. In Theorem 12.4 it is shown that every τ-transition is matched by a reduction. Conversely, for every reduction there exists a corresponding τ-transition, as is proved in Theorem 14.4.

Theorem 12.

1. *If* $P \xrightarrow{\bar{x}y} P'$ *then* $P \equiv \bar{x}y\,P'$.
2. *If* $P \xrightarrow{x(y)} P'$ *then* $P \equiv x(y)\,P'$.
3. *If* $P \xrightarrow{\sigma_y\,\bar{x}y} P'$ *then* $P \equiv \sigma_y\,\bar{x}y\,P'$.

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[2] $\tau\,Q$ and $\tau\,Q'$ are alpha-convertible if Q and Q' are.

4. If $P \xrightarrow{\tau} P'$ then $P \to P'$.

Proof. Induction on the proofs exploiting Proposition 7. □

The proof of Theorem 14 relies on the following lemma. This lemma is the main technical result of the paper.

Lemma 13. *Let $P \equiv Q$.*

1. If $P \xrightarrow{\alpha} P'$ then $Q \xrightarrow{\alpha} Q'$ for some Q' such that $P' \equiv Q'$.
2. If $Q \xrightarrow{\alpha} Q'$ then $P \xrightarrow{\alpha} P'$ for some P' such that $P' \equiv Q'$.

Proof. We prove this lemma by induction on the proofs of $P \xrightarrow{\alpha} P'$, $Q \xrightarrow{\alpha} Q'$, and $P \equiv Q$. We only consider proofs of $P \xrightarrow{\alpha} P'$ and $Q \xrightarrow{\alpha} Q'$ of minimal complexity. The complexity of a proof is determined by those nodes in the proof where the rule 1. is applied. The more this rule is applied towards the root of the proof, the smaller its complexity is. In the proof we exploit Proposition 4, 5, 10, and 11. □

We conclude this section with

Theorem 14.

1. If $P \equiv \bar{x}y\, P'$ then $P \xrightarrow{\bar{x}y} P''$ for some P'' such that $P'' \equiv P'$.
2. If $P \equiv x(y)\, P'$ then $P \xrightarrow{x(y)} P''$ for some P'' such that $P'' \equiv P'$.
3. If $P \equiv \sigma_y\, \bar{x}y\, P'$ then $P \xrightarrow{\sigma_y\, \bar{x}y} P''$ for some P'' such that $P'' \equiv P'$.
4. If $P \to P'$ then $P \xrightarrow{\tau} P''$ for some P'' such that $P'' \equiv P'$.

Proof. From Lemma 13 we conclude 1., 2., and 3. The fourth case we prove by induction on the proof of $P \to P'$. □

5 Basic π_ϵI-calculus

In this section we restrict our attention to a subcalculus of π_ϵ-calculus which only gives rise to *internal mobility* (see [San96]) called π_ϵI-calculus. The reduction system of Section 2 is easily adapted. Like for ordinary π-calculus, the labelled transition system for the subcalculus is much simpler than the one for the full calculus given in Section 3. The relation between the two systems is similar to the one presented in Section 4.

The subcalculus is obtained by restricting the set of particles. We do not consider *free* outputs $\bar{x}y$ but only *bound* ones $(\nu y)\,\bar{x}y$, from now on abbreviated to $\bar{x}(y)$.

Definition 15. The set of *particles* is given by

$$\pi ::= \bar{x}(y) \mid x(y) \mid (\nu x)$$

The particle $\bar{x}(y)$ is a binder with $\operatorname{bn}(\bar{x}(y)) = \{y\}$ and $\operatorname{fn}(\bar{x}(y)) = \{x\}$. The structural congruence $\equiv$ is defined by all the rules of Definition 2 but the rule 5. The latter rule can be derived from the other ones (see [Bre97]). The reduction relation $\rightarrow$ is obtained from Definition 3 by replacing the axiom 1. by

1. $x(y)\,P \mid \bar{x}(y)\,Q \rightarrow (\nu y)\,(P \mid Q)$

Note that we only encounter alpha-conversion and no substitution in the reduction system for π_ϵI-calculus.

In the labelled transition system we do not need the extended scopes of Definition 6 we used in Section 3.

Definition 16. The set of *actions* is given by

$$\alpha ::= \bar{x}(y) \mid x(y) \mid \tau$$

The transition relation is presented next. We have omitted the symmetric versions of the rules 7. and 8.

Definition 17. The *transition relation* $\rightarrow$ is defined as the smallest labelled relation over processes satisfying

1. $\dfrac{P' \xrightarrow{\alpha} Q'}{P \xrightarrow{\alpha} Q}$ P and P', and Q and Q' are alpha-convertible

2. $\bar{x}(y)\,P \xrightarrow{\bar{x}(y)} P$

3. $x(y)\,P \xrightarrow{x(y)} P$

4. $\dfrac{P \xrightarrow{\alpha} P'}{\pi\,P \xrightarrow{\alpha} \pi\,P'}$ $\operatorname{n}(\alpha) \cap \operatorname{bn}(\pi) = \emptyset$ and $\operatorname{n}(\pi) \cap \operatorname{bn}(\alpha) = \emptyset$

5. $\dfrac{P \xrightarrow{x(z)} P'}{\bar{x}(y)\,P \xrightarrow{\tau} (\nu y)\,(P'[y\!/\!z])}$ $x \neq y$

6. $\dfrac{P \xrightarrow{\bar{x}(z)} P'}{x(y)\,P \xrightarrow{\tau} (\nu y)\,(P'[y\!/\!z])}$ $x \neq y$

7. $\dfrac{P \xrightarrow{\alpha} P'}{P \mid Q \xrightarrow{\alpha} P' \mid Q}$ $\operatorname{bn}(\alpha) \cap \operatorname{fn}(Q) = \emptyset$

8. $\dfrac{P \xrightarrow{x(y)} P' \quad Q \xrightarrow{\bar{x}(y)} Q'}{P \mid Q \xrightarrow{\tau} (\nu y)\,(P' \mid Q')}$

Some remarks:

– The rules 1., 4., and 7., and the axiom 3. correspond to the rules 1., 4., and 9., and the axiom 3. of Definition 9.

- The axiom 2. and the rules 5., 6., and 8. are the obvious modifications of the axiom 2., and the rules 7., 8., and 10. of Definition 9.
- Note that we do use substitution in the rules 5. and 6. In the transition corresponding to the reduction

$$\bar{x}(y)\, x(z)\, z(w)\, \bar{y}(w)\, 0 \rightarrow (\nu y)\, y(w)\, \bar{y}(w)\, 0 \tag{6}$$

(a proof of this reduction is given in [Bre97]) z in $z(w)$ and y in $\bar{y}(w)$ are identified:

$$\bar{x}(y)\, x(z)\, z(w)\, \bar{y}(w)\, 0 \xrightarrow{\tau} (\nu y)\, y(w)\, \bar{y}(w)\, 0.$$

This identification cannot be brought about by alpha-conversion of the process $\bar{x}(y)\, x(z)\, z(w)\, \bar{y}(w)\, 0$.

We conclude this section with a correspondence theorem.

Theorem 18.

1. *If $P \xrightarrow{\bar{x}(y)} P'$ then $P \equiv \bar{x}(y)\, P'$.*

2. *If $P \xrightarrow{x(y)} P'$ then $P \equiv x(y)\, P'$.*

3. *If $P \xrightarrow{\tau} P'$ then $P \rightarrow P'$.*

4. *If $P \equiv \bar{x}(y)\, P'$ then $P \xrightarrow{\bar{x}(y)} P''$ for some P'' such that $P'' \equiv P'$.*

5. *If $P \equiv x(y)\, P'$ then $P \xrightarrow{x(y)} P''$ for some P'' such that $P'' \equiv P'$.*

6. *If $P \rightarrow P'$ then $P \xrightarrow{\tau} P''$ for some P'' such that $P'' \equiv P'$.*

Proof. Similar to the proof of Theorem 12 and 14. □

6 Related work

The only three other papers which discuss the relation between a reduction system and a labelled transition system we are aware of are Ferrari's [Fer96], Milner's [Mil92], and Honda and Yoshida's [HY93]. Ferrari considers a CCS-like calculus. Since the calculus contains no binders, the problem of relating a reduction system and a labelled transition system becomes much simpler. Milner, and Honda and Yoshida focus on π-calculus. They both use the rule

$$1.'\quad \frac{P' \xrightarrow{\alpha} Q'}{P \xrightarrow{\alpha} Q} \quad P \equiv P' \text{ and } Q \equiv Q'$$

instead of the rule 1. of Definition 9. In their setting Lemma 13, the main technical result of this paper, becomes trivial. Their rule is less structural than ours. Furthermore, the rule 1. can easily be distributed over the other axioms and rules (compare the labelled transition system of Milner et al. [MPW92, page 46] and the one of Sangiorgi [San92, page 30]). This is not the case for the other rule.

In the conclusion of [MP95], Montanari and Pistore consider relaxing the sequencing power of prefixing. Instead of a reduction system or a labelled tran-

sition system, they use a graph rewriting system. In their setting, enablement can easily be accommodated (as long as one does not consider replication).

Conclusion

From our case study we can conclude that the problem of reconstructing a labelled transition system from a reduction system is far from easy. Although the reduction system for π_ϵ-calculus is rather close to the one for ordinary π-calculus, we encounter in the labelled transition system for π_ϵ-calculus extended scopes and various new rules.

In [Mil93, page 37], Milner first presented π_ϵ-calculus with enablement as its new feature. The fact that π_ϵ-calculus has self communication was already observed by Bellin and Scott [BS94, page 15]. But the presence of extended scope extrusion in π_ϵ-calculus—although maybe not very surprising—only occurred to us when we developed the labelled transition system.

The labelled transition system for π_ϵ-calculus might be the basis for the development of a (possibly fully abstract with respect to some form of bisimulation) denotational semantics for the calculus. Here we can make fruitful use of the work of Fiore, Moggi, and Sangiorgi [FMS96], Hennessy [Hen96], and Stark [Sta96].

The labelled transition system for π_ϵI-calculus, the subcalculus with only internal mobility, is much simpler than the one for the full calculus. This provides another indication that external mobility is responsible for much of the semantic complications (cf. [San96]).

Although we only consider internal mobility in π_ϵI-calculus, we do use substitution in the labelled transition system. In πI-calculus only alpha-conversion is needed (see [San96, Section 2.2]). This suggests that the absence of substitution in πI-calculus is just a property of the calculus, rather than a consequence of its restriction to internal mobility. Whether the substitutions used in π_ϵI-calculus are of a special kind (the substituted name is always bound by a generated restriction) needs further study.

Another topic reserved for later treatment is the study of bisimulation. The definitions of barbed, early, ground, late, and open bisimulation for π-calculus can be adapted straightforwardly to our setting (see [Bre97]). We are interested in the connection with bisimulation for action structures (for π_ϵ-calculus) given by Milner in [Mil93].

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