



On Contact Graphs of Paths on a Grid

Zakir Deniz¹, Esther Galby²(✉), Andrea Munaro², and Bernard Ries²

¹ Department of Mathematics, Duzce University, Duzce, Turkey
zakirdeniz@duzce.edu.tr

² Department of Informatics, Decision Support & Operations Research,
University of Fribourg, Fribourg, Switzerland
{esther.galby, andrea.munaro, bernard.ries}@unifr.ch

Abstract. In this paper we consider *Contact graphs of Paths on a Grid* (*CPG graphs*), i.e. graphs for which there exists a family of interiorly disjoint paths on a grid in one-to-one correspondence with their vertex set such that two vertices are adjacent if and only if the corresponding paths touch at a grid-point. Our class generalizes the well studied class of VCPG graphs (see [1]). We examine CPG graphs from a structural point of view which leads to constant upper bounds on the clique number and the chromatic number. Moreover, we investigate the recognition and 3-colorability problems for B_0 -CPG, a subclass of CPG. We further show that CPG graphs are not necessarily planar and not all planar graphs are CPG.

1 Introduction

Asinowski et al. [3] introduced the class of *vertex intersection graphs of paths on a grid*, referred to as *VPG graphs*. An undirected graph $G = (V, E)$ is called a *VPG graph* if one can associate a path on a grid with each vertex such that two vertices are adjacent if and only if the corresponding paths intersect on at least one grid-point. It is not difficult to see that the class of VPG graphs coincides with the class of string graphs, i.e. intersection graphs of curves in the plane (see [3]).

A natural restriction which was forthwith considered consists in limiting the number of *bends* (i.e. 90° turns at a grid-point) that the paths may have: an undirected graph $G = (V, E)$ is a B_k -*VPG graph*, for some integer $k \geq 0$, if one can associate a path on a grid having at most k bends with each vertex such that two vertices are adjacent if and only if the corresponding paths intersect on at least one grid-point. Since their introduction, B_k -VPG have been extensively studied (see [2, 3, 5, 7–9, 14, 15, 18–20]).

A notion closely related to intersection graphs is that of *contact graphs*. Such graphs can be seen as a special type of intersection graphs of geometrical objects in which these objects are not allowed to have common interior points but only to touch each other. Contact graphs of various types of objects have been studied in the literature (see, e.g., [1, 10, 11, 21–23]). In this paper, we consider *Contact graphs of Paths on a Grid* (*CPG graphs* for short) which are defined as follows.



Fig. 1. Examples of types of contact between two paths (the endpoints of a path are marked by an arrow).

A graph G is a *CPG graph* if the vertices of G can be represented by a family of interiorly disjoint paths on a grid, two vertices being adjacent in G if and only if the corresponding paths touch, i.e. share a grid-point which is an endpoint of at least one of the two paths (see Fig. 1). Note that this class is hereditary, i.e. closed under vertex deletion. Similarly to VPG, a B_k -CPG graph is a CPG graph admitting a representation in which each path has at most k bends. Clearly, any B_k -CPG graph is also a B_k -VPG graph.

Aerts and Felsner [1] considered a similar family of graphs, namely those admitting a *Vertex Contact representation of Paths on a Grid (VCPG for short)*. The vertices of such graphs can be represented by a family of interiorly disjoint paths on a grid, but the adjacencies are defined slightly differently: two vertices are adjacent if and only if the endpoint of one of the corresponding paths touches an interior point of the other corresponding path (observe that this is equivalent to adding the constraint forbidding two paths from having a common endpoint, i.e. contacts as in Fig. 1a on the right). This class has been considered by other authors as well (see [6, 7, 14, 19, 24]).

It is not difficult to see that graphs admitting a VCPG are planar (see [1]) and it immediately follows from the definition that those graphs are CPG graphs. This containment is in fact strict even when restricted to planar CPG graphs, as there exist, in addition to nonplanar CPG graphs, planar graphs which are CPG but do not admit a VCPG.

To the best of our knowledge, the class of CPG graphs has never been studied in itself and our present intention is to provide some structural properties (see Sect. 3). By considering a specific weight function on the vertices, we provide upper bounds on the number of edges in CPG graphs as well as on the clique number and the chromatic number (see Sect. 3). In particular, we show that B_0 -CPG graphs are 4-colorable and that 3-COLORABILITY restricted to B_0 -CPG is NP-complete (see Sect. 5). We further prove that recognizing B_0 -CPG graphs is NP-complete. Additionally, we show that the classes of CPG graphs and planar graphs are incomparable (see Sect. 4).

2 Preliminaries

Throughout this paper, all considered graphs are undirected, finite and simple. For any graph theoretical notion not defined here, we refer the reader to [13].

Let $G = (V, E)$ be a graph with vertex set V and edge set E . The *degree* of a vertex $v \in V$, denoted by $d(v)$, is the number of neighbors of v in G . A graph G is k -*regular* if the degree of every vertex in G is $k \geq 0$. A *clique* (resp. *stable set*) in G is a set of pairwise adjacent (resp. nonadjacent) vertices. The graph obtained from G by deleting a vertex $v \in V$ is denoted by $G - v$. For a given graph H , G is H -*free* if it contains no induced subgraph isomorphic to H .

As usual, K_n (resp. C_n) denotes the complete graph (resp. chordless cycle) on n vertices and $K_{m,n}$ denotes the complete bipartite graph with bipartition (V_1, V_2) such that $|V_1| = m$ and $|V_2| = n$. Given a graph G , the *line graph* of G , denoted by $L(G)$, is the graph such that each vertex v_e in $L(G)$ corresponds to an edge e in G and two vertices are adjacent in $L(G)$ if and only if their corresponding edges in G have a common endvertex.

A graph G is *planar* if it can be drawn in the plane without crossing edges; such a drawing is then called a *planar embedding* of G . A planar embedding divides the plane into several regions referred to as *faces*. A planar graph is *maximally planar* if adding any edge renders it nonplanar. A maximally planar graph has exactly $2n - 4$ faces, where n is the number of vertices in the graph. A graph H is a *minor* of a graph G , if H can be obtained from G by deleting edges and vertices and by contracting edges. It is well-known that a graph is planar if and only if it does not contain K_5 or $K_{3,3}$ as a minor [13].

A *coloring* of a graph G is a mapping \mathbf{c} associating with every vertex u an integer $\mathbf{c}(u)$, called a *color*, such that $\mathbf{c}(v) \neq \mathbf{c}(u)$ for every edge uv . If at most k distinct colors are used, \mathbf{c} is called a k -*coloring*. The smallest integer k such that G admits a k -coloring is called the *chromatic number* of G , denoted by $\chi(G)$.

Consider a rectangular grid \mathcal{G} where the horizontal lines are referred to as *rows* and the vertical lines as *columns*. The grid-point lying on row x and column y is denoted by (x, y) . An *interior point* of a path P on \mathcal{G} is a point belonging to P and different from its endpoints; the *interior* of P is the set of all its interior points. A graph $G = (V, E)$ is *CPG* if there exists a collection \mathcal{P} of interiorly disjoint paths on a grid \mathcal{G} such that \mathcal{P} is in one-to-one correspondence with V and two vertices are adjacent in G if and only if the corresponding paths touch; if every path in \mathcal{P} has at most k bends, G is B_k -CPG. The pair $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ is a *CPG representation* of G , and more specifically a k -*bend CPG representation* if every path in \mathcal{P} has at most k bends. In the following, the path representing some vertex u in a CPG representation \mathcal{R} of a graph G is denoted by $P_u^{\mathcal{R}}$, or simply P_u if it is clear from the context.

Let $G = (V, E)$ be a CPG graph and $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ be a CPG representation of G . A grid-point p is of *type I* if it corresponds to an endpoint of four paths in \mathcal{P} (see Fig. 2a), and of *type II* if it corresponds to an endpoint of two paths in \mathcal{P} and an interior point of a third path in \mathcal{P} (see Fig. 2b).

For any grid-point p , we denote by $\tau(p)$ the number of edges in the subgraph induced by the vertices whose corresponding paths contain or have p as an endpoint. Note that this subgraph is a clique and so $\tau(p) = \binom{j}{2}$ if j paths touch at grid-point p .

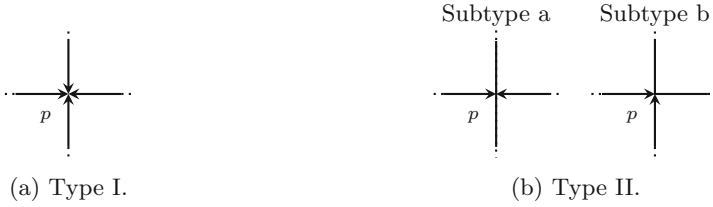


Fig. 2. Two types of grid-points.

For any path P , we denote by $\overset{\circ}{P}$ (resp. $\partial(P)$) the interior (resp. endpoints) of P . For a vertex $u \in V$, we define the *weight of u with respect to \mathcal{R}* , denoted by $w_u^{\mathcal{R}}$ or simply w_u if it is clear from the context, as follows. Let q_u^i ($i = 1, 2$) be the endpoints of the corresponding path P_u in \mathcal{P} and consider, for $i = 1, 2$,

$$w_u^i = |\{P \in \mathcal{P} \mid q_u^i \in \overset{\circ}{P}\}| + \frac{1}{2} \cdot |\{P \in \mathcal{P} \mid P \neq P_u \text{ and } q_u^i \in \partial(P)\}|.$$

Then $w_u = w_u^1 + w_u^2$.

Observation 1. *Let $G = (V, E)$ be a CPG graph and $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ be a CPG representation of G . For any vertex $u \in V$ and $i = 1, 2$, $w_u^i \leq \frac{3}{2}$ where equality holds if and only if q_u^i is a grid-point of type I or II.*

Indeed, the contribution of q_u^i to w_u^i is maximal if all four grid-edges containing q_u^i are used by paths of \mathcal{P} , which may only happen when q_u^i is a grid-point of type I or II.

Remark. In fact, we have $w_u^i \in \{0, \frac{1}{2}, 1, \frac{3}{2}\}$ for any vertex $u \in V$ and $i = 1, 2$.

Observation 2. *Let $G = (V, E)$ be a CPG graph and $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ be a CPG representation of G . Then*

$$|E| \leq \sum_{u \in V} w_u,$$

where equality holds if and only if all paths of \mathcal{P} pairwise touch at most once.

Indeed, if $uv \in E$, we may assume that either an endpoint of P_u touches the interior of P_v , or P_u and P_v have a common endpoint. In the first case, the edge uv is fully accounted for in the weight of u , and in the second case, the edge uv is accounted for in both w_u and w_v by one half. The characterization of equality then easily follows.

3 Structural Properties of CPG Graphs

In this section, we investigate CPG graphs from a structural point of view and present some useful properties which we will further exploit.

Lemma 1. *A CPG graph is either 6-regular or has a vertex of degree at most 5.*

Proof. If $G = (V, E)$ is a CPG graph and \mathcal{R} is a CPG representation of G , by combining Observations 1 and 2, we obtain

$$\sum_{u \in V} d(u) = 2|E| \leq 2 \sum_{u \in V} w_u \leq 2 \sum_{u \in V} \left(\frac{3}{2} + \frac{3}{2}\right) = 6|V|. \quad \square$$

Remark. We can show that there exists an infinite family of 6-regular CPG graphs. Due to lack of space, this proof is here omitted but can be found in the full version [12].

For B_1 -CPG graphs, we can strengthen Lemma 1 as follows.

Proposition 1. *Every B_1 -CPG graph has a vertex of degree at most 5.*

Proof. Let $G = (V, E)$ be a B_1 -CPG graph and \mathcal{R} be a 1-bend CPG representation of G . Denote by p the upper-most endpoint of a path among the left-most endpoints in \mathcal{R} , and by P_x (with $x \in V$) an arbitrary path having p as an endpoint. Since \mathcal{R} is a 1-bend CPG representation, no path uses the grid-edge on the left of p , for otherwise p would not be a left-most endpoint. Therefore, p contributes to the weight of x with respect to \mathcal{R} by at most 1 and, by Observations 1 and 2, we have

$$\sum_{u \in V} d(u) = 2|E| \leq 2(w_x + \sum_{u \neq x} w_u) \leq 6|V| - 1,$$

which implies the existence of a vertex of degree at most 5. □

A natural question that arises when considering CPG graphs is whether they may contain large cliques. It immediately follows from Observation 2 that CPG graphs cannot contain K_n , for $n \geq 8$. This can be further improved as shown in the next result.

Theorem 1. *CPG graphs are K_7 -free.*

Proof. Since the class of CPG graphs is hereditary, it is sufficient to show that K_7 is not a CPG graph. Suppose, to the contrary, that K_7 is a CPG graph and consider a CPG representation $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ of K_7 . Observe first that the weight of every vertex with respect to \mathcal{R} must be exactly $2 \cdot 3/2$, as otherwise by Observation 1, we would have $\sum_{u \in V} w_u < 3|V| = 21 = |E|$ which contradicts Observation 2. This implies in particular that every grid-point corresponding to an endpoint of a path is either of type I or II. Furthermore, any two paths must touch at most once, for otherwise by Observation 2, $|E| < \sum_{u \in V} w_u = 3|V| = |E|$. Hence, if we denote by P_I (resp. P_{II}) the set of grid-points of type I (resp. type II), then since $\tau(p) = 6$ for all $p \in P_I$ and $\tau(p) = 3$ for all $p \in P_{II}$, we have that $6|P_I| + 3|P_{II}| = 21$, which implies $|P_{II}| \neq 0$. Suppose that there exists a path P_u having one endpoint corresponding to a grid-point of type I and the

other corresponding to a grid-point of type II. Since the corresponding vertex u has degree 6, P_u must then properly contain an endpoint of another path which, as first observed, necessarily corresponds to a grid-point of type II. But vertex u would then have degree $3 + 2 + 2$ as no two paths touch more than once, a contradiction. Hence, every path has both its endpoints of the same type. But then, $|P_I| = 0$; indeed, if there exists a path having both its endpoints of type I, since no two paths touch more than once, this implies that every path has both its endpoints of type I, i.e. $|P_{II}| = 0$, a contradiction. Now, if we consider each grid-point of type II as a vertex and connect any two such vertices when the corresponding grid-points belong to a same path, then we obtain a planar embedding of a 4-regular graph on 7 vertices. But this contradicts the fact that every 4-regular graph on 7 vertices contains $K_{3,3}$ as a minor (a proof of this result can be found in the full version [12]). \square

However, CPG graphs may contain cliques on 6 vertices as shown in Proposition 2. Due to lack of space, its proof is omitted here and can be found in the full version [12].

Proposition 2. K_6 is in B_2 -CPG $\setminus B_1$ -CPG.

We conclude this section with a complexity result pointing towards the fact that there may not be a polynomial characterization of B_0 -CPG graphs. Let us first introduce rectilinear planar graphs: a graph G is *rectilinear planar* if it admits a rectilinear planar drawing, i.e. a drawing mapping each edge to a horizontal or vertical segment.

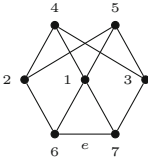
Theorem 2. RECOGNITION is NP-complete for B_0 -CPG graphs.

Proof. We show that a graph G is rectilinear planar if and only if its line graph $L(G)$ is B_0 -CPG. As RECOGNITION for rectilinear planar graphs was shown to be NP-complete in [17], this concludes the proof. Suppose G is a rectilinear planar graph and let \mathcal{D} be the collection of horizontal and vertical segments in a rectilinear planar drawing of G . It is not difficult to see that the contact graph of \mathcal{D} is isomorphic to $L(G)$. Conversely, assume that $L(G)$ is a B_0 -CPG graph and consider a 0-bend CPG representation $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ of $L(G)$. Since $L(G)$ is $K_{1,3}$ -free [4], every path in \mathcal{P} has at most two contact points. Thus, by eventually shortening paths, we may assume that contacts only happen at endpoints of paths. Therefore, \mathcal{R} induces a rectilinear planar drawing of G , where each vertex corresponds to a contact point in \mathcal{R} and each edge is mapped to its corresponding path in \mathcal{P} . \square

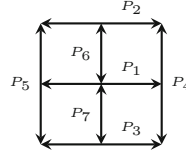
4 Planar CPG Graphs

In this section, we focus on planar graphs and their relation with CPG graphs. In particular, we show that not every planar graph is CPG and not all CPG graphs are planar.¹

¹ We can further show that not all CPG graphs are 1-planar as $K_7 - E(K_3)$ is CPG but not 1-planar [25].



(a) A nonplanar graph G .



(b) A 0-bend CPG representation of G .

Fig. 3. A B_0 -CPG graph containing $K_{3,3}$ as a minor (contract the edge e).

Lemma 2. *If G is a CPG graph for which there exists a CPG representation containing no grid-point of type I or II.a, then G is planar. In particular, if G is a triangle-free CPG graph, then G is planar.*

Proof. Let $G = (V, E)$ be a CPG graph for which there exists a CPG representation \mathcal{R} containing no grid-point of type I or II.a. By considering each path of \mathcal{R} as a curve in the plane, it follows that G is a curve contact graph having a representation (namely \mathcal{R}) in which any point in the plane belongs to at most three curves. Furthermore, whenever a point in the plane belongs to the interior of a curve \mathcal{C} and corresponds to an endpoint of two other curves, then those two curves lie on the same side of \mathcal{C} (recall that there is no grid-point of type II.a). Hence, it follows from Proposition 2.1 in [21] that G is planar.

If G is a triangle-free CPG graph, then no CPG representation of G contains grid-points of type I or II.a. Hence, G is planar. □

Remark. Since $K_{3,3}$ is a triangle-free nonplanar graph, it follows from Lemma 2 that $K_{3,3}$ is not CPG. Therefore, CPG graphs are $K_{3,3}$ -free. Observe however that for any $k \geq 0$, B_k -CPG is not a subclass of planar graphs as there exist B_0 -CPG graphs which are not planar (see Fig. 3).

It immediately follows from [7] that all triangle-free planar graphs are B_1 -CPG; hence, we have the following corollary.

Corollary 1. *If a graph G is triangle-free, then G is planar if and only if G is B_1 -CPG.*

The next result allows us to detect planar graphs that are not CPG.

Lemma 3. *Let $G = (V, E)$ be a planar graph. If G is a CPG graph, then G has at most $4|V| - 2f + 4$ vertices of degree at most 3, where f denotes the number of faces of G . In particular, if G is maximally planar, then G has at most 12 vertices of degree at most 3.*

Proof. Let $G = (V, E)$ be a planar CPG graph and $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ a CPG representation of G . Denote by U the subset of vertices in G of degree at most 3. If a path P_u , with $u \in U$, touches every other path in \mathcal{P} at most once, then, since at least one endpoint of P_u is then not a grid-point of type I or II, the weight of

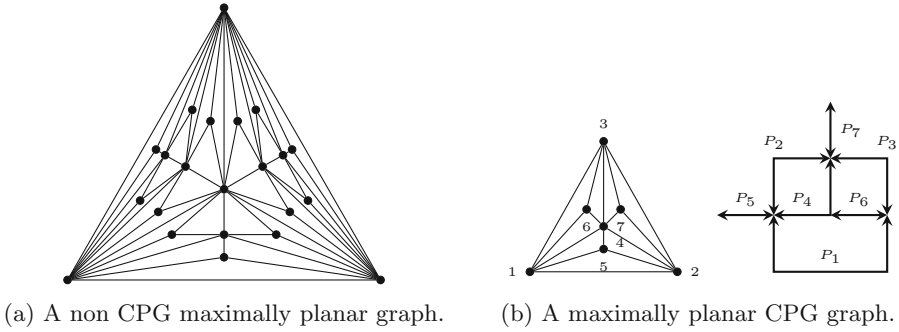


Fig. 4. Two maximally planar graphs.

u with respect to \mathcal{R} is at most $3/2 + 1$. Thus, if we assume that this is the case for all paths whose corresponding vertex is in U , we have by Observation 2

$$|E| \leq \left(\frac{3}{2} + 1\right)|U| + 3(|V| - |U|) = 3|V| - \frac{|U|}{2}.$$

On the other hand, if there exists $u \in U$ such that P_u touches some path more than once, then the above inequality still holds as the corresponding edge is already accounted for. Using the fact that $f = |E| - |V| + 2$ (Euler’s formula), we obtain the desired upper bound. Moreover, if G is maximally planar, then $f = 2|V| - 4$ and so $|U| \leq 12$. \square

Remark. In Fig. 4a, we give an example of a maximally planar graph which is not CPG due to Lemma 3. It is constructed by iteratively adding a vertex in a triangular face, starting from the triangle, so that it has exactly 13 vertices of degree 3. There exist however maximally planar graphs which are CPG (see Fig. 4b). Note that maximally planar graphs do not admit a VCPG [1].

5 Coloring CPG Graphs

In this section, we provide tight upper bounds on the chromatic number of B_k -CPG graphs for different values of k and investigate the 3-COLORABILITY problem for CPG graphs. The proof of the following result is an easy exercise left to the reader (see the full version [12]).

Theorem 3. *CPG graphs are 6-colorable.*

Remark. Since K_6 is B_2 -CPG, this bound is tight for B_k -CPG graphs with $k \geq 2$. We leave as an open problem whether this bound is also tight for B_1 -CPG graphs (note that it is at least 5 since K_5 is B_1 -CPG).

Theorem 4. *B_0 -CPG graphs are 4-colorable. Moreover, K_4 is a 4-chromatic B_0 -CPG graph.*

Proof. Let G be a B_0 -CPG graph and $\mathcal{R} = (\mathcal{G}, \mathcal{P})$ a 0-bend CPG representation of G . Denote by \mathcal{L} (resp. \mathcal{C}) the set of rows (resp. columns) of \mathcal{G} on which lies at least one path of \mathcal{P} . Since the representation contains no bend, if A is a row in \mathcal{L} (resp. column in \mathcal{C}), then the set of vertices having their corresponding path on A induces a collection of disjoint paths in G . If $B \neq A$ is another row in \mathcal{L} (resp. column in \mathcal{C}), then no path in A touches a path in B . Hence, it suffices to use two colors to color the vertices having their corresponding path in a row of \mathcal{L} and two other colors to color the vertices having their corresponding path in a column of \mathcal{C} to obtain a proper coloring of G . \square

It immediately follows from a result in [22] that the 3-COLORABILITY problem is NP-complete in CPG, even if the graph admits a representation in which each grid-point belongs to at most two paths. We conclude this section by a strengthening of this result.

Theorem 5. 3-COLORABILITY is NP-complete in B_0 -CPG.

Proof. We exhibit a polynomial reduction from 3-COLORABILITY restricted to planar graphs of maximum degree 4, which was shown to be NP-complete in [16].

Let $G = (V, E)$ be a planar graph of maximum degree 4. It follows from [26] that G admits a grid embedding where each vertex is mapped to a grid-point and each edge is mapped to a grid-path with at most 4 bends, in such a way that all paths are interiorly disjoint (such an embedding can be obtained in linear time). Denote by $\mathcal{D} = (\mathcal{V}, \mathcal{E})$ such an embedding, where \mathcal{V} is the set of grid-points in one-to-one correspondence with V and \mathcal{E} is the set of grid-paths in one-to-one correspondence with E . For any vertex $u \in V$, we denote by (x_u, y_u) the grid-point in \mathcal{V} corresponding to u and by P_u^N (resp. P_u^S) the path of \mathcal{E} , if any, having (x_u, y_u) as an endpoint and using the grid-edge above (resp. below) (x_u, y_u) . For any edge $e \in E$, we denote by P_e the path in \mathcal{E} corresponding to e . We construct from \mathcal{D} a 0-bend CPG representation \mathcal{R} in such a way that the corresponding graph G' is 3-colorable if and only if G is 3-colorable.

By eventually adding rows and columns to the grid, we may assume that the interior of each path P in \mathcal{E} is surrounded by an empty region, i.e. no path $P' \neq P$ or grid-point of \mathcal{V} lies in the interior of this region. In the following, we denote this region by \mathcal{R}_P (delimited by red dashed lines in every subsequent figure) and assume, without loss of generality, that it is always large enough for the following operations.

We first associate with every vertex $u \in V$ a vertical path P_u containing the grid-point (x_u, y_u) as follows. If P_u^N (resp. P_u^S) is not defined, the top (resp. lower) endpoint of P_u is $(x_u, y_u + \varepsilon)$ (resp. $(x_u, y_u - \varepsilon)$) for a small enough ε so that the segment $[(x_u, y_u), (x_u, y_u + \varepsilon)]$ (resp. $[(x_u, y_u), (x_u, y_u - \varepsilon)]$) touches no path of \mathcal{E} . If P_u^N has at least one bend, then the top endpoint of P_u lies at the border of $\mathcal{R}_{P_u^N}$ on column x_u (see Fig. 5a). If P_u^N has no bend, then the top endpoint of P_u lies at the middle of P_u^N (see Fig. 5b). Similarly, we define the lower endpoint of P_u according to P_u^S : if P_u^S has at least one bend, then the lower endpoint of P_u lies at the border of $\mathcal{R}_{P_u^S}$ on column x_u , otherwise it lies at the middle of P_u^S .

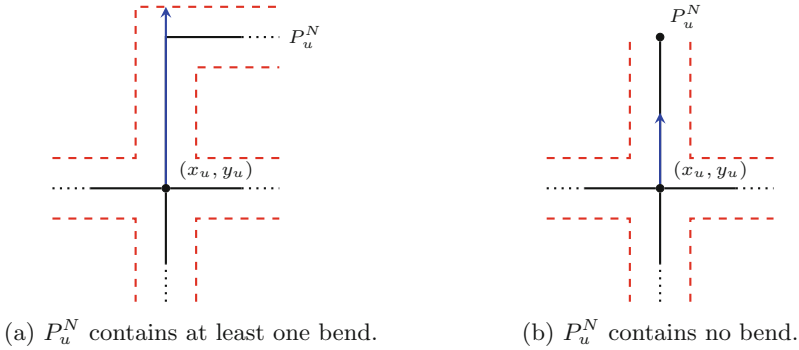


Fig. 5. Constructing the path P_u corresponding to vertex u (in blue). (Color figure online)

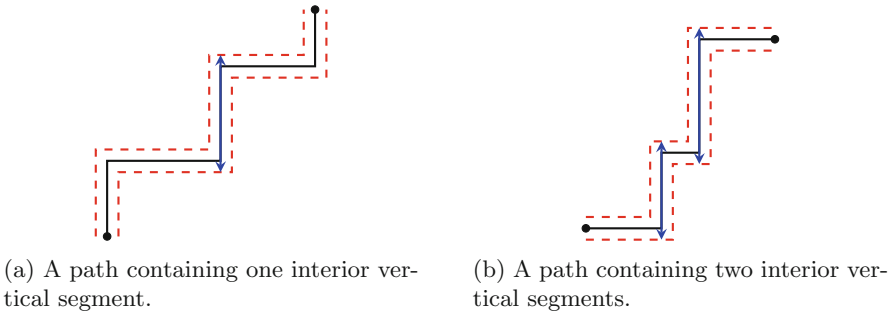


Fig. 6. Replacing interior vertical segments by 0-bend paths (in blue). (Color figure online)

For any path P of \mathcal{E} with at least two bends, an *interior vertical segment* of P is a vertical segment of P containing none of its endpoints (note that since every path in \mathcal{E} has at most 4 bends, it may contain at most two interior vertical segments). We next replace every interior segment of P by a slightly longer vertical path touching the border of \mathcal{R}_P (see Fig. 6).

We finally introduce two gadgets H (see Fig. 7) and H' , where H' is the subgraph of H induced by $\{b, c, 4, 5, 6, 7, 8, 9, 10\}$, as follows. Denote by \mathcal{P}' the set of vertical paths introduced so far and by V' the set of vertices of the contact graph of \mathcal{P}' . Observe that V' contains a copy of V and that two vertices are adjacent in the contact graph of \mathcal{P}' if and only if they are both copies of vertices in V and the path P of \mathcal{P} corresponding to the edge between these two copies is a vertical path with no bend. Now, along each path P_{uv} of \mathcal{P} such that the vertical paths P_u and P_v of \mathcal{P}' do not touch, we add gadgets H and H' as follows. Let P_1, \dots, P_k be the vertical paths of \mathcal{P}' encountered in order when going along P_{uv} from (x_u, y_u) to (x_v, y_v) and let u_j be the vertex of V' corresponding to P_j , for $1 \leq j \leq k$. Note that P_1 (resp. P_k) is the path corresponding to vertex

$u = u_1$ (resp. $v = u_k$) and that P_j , for $2 \leq j \leq k - 1$, is a path corresponding to an interior vertical segment of P_{uv} (this implies in particular that $k \leq 4$). We add the gadget H' in between u_1 and u_2 by identifying u_1 with b and u_2 with c . Moreover, for any $2 \leq j \leq k - 1$, we add the gadget H in between u_j and u_{j+1} by identifying u_j with b and u_{j+1} with a (see Fig. 8 where $k = 4$ and each box labeled H (resp. H') means that gadget H (resp. H') has been added by identifying the vertex lying to the left of the box to b and the vertex lying on the right of the box to a (resp. c)).

The resulting graph G' remains B_0 -CPG. Indeed, we may add 0-bend CPG representations of the gadgets H and H' inside $\mathcal{R}_{P_{uv}}$ and at different heights so that they do not touch any other such gadget, as shown in Fig. 9. In the full version [12], we give a local example of the resulting 0-bend CPG representation \mathcal{R} .

We now show that G is 3-colorable if and only if G' is. To this end, we prove the following.

Claim 1

- In any 3-coloring \mathbf{c} of H' , we have $\mathbf{c}(b) \neq \mathbf{c}(c)$.
- In any 3-coloring \mathbf{c} of H , we have $\mathbf{c}(a) = \mathbf{c}(b)$ and $\mathbf{c}(b) \neq \mathbf{c}(c)$.

Proof. Let $\mathbf{c}: \{a, b, c, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10\} \rightarrow \{blue, red, green\}$ be a 3-coloring of H and assume without loss of generality that $\mathbf{c}(b) = blue$. Clearly, at least two vertices among 4, 6 and 8 have the same color. If vertices 4, 6 and 8 all have the same color, say *red*, then either $\mathbf{c}(7) = blue$ and $\mathbf{c}(9) = green$, or $\mathbf{c}(7) = green$ and $\mathbf{c}(9) = blue$. Therefore, $\{\mathbf{c}(5), \mathbf{c}(10)\} = \{blue, green\}$ and since c is adjacent to all three colors, we then obtain a contradiction. Now if vertices 4 and 8 have the same color, say *red*, then vertex 6 has color *green* and both 7 and 9 have color *blue*, a contradiction. Hence, either $\mathbf{c}(4) = \mathbf{c}(6) \neq \mathbf{c}(8)$ or $\mathbf{c}(8) = \mathbf{c}(6) \neq \mathbf{c}(4)$. By symmetry, we may assume that vertices 4 and 6 have the same color, say *red*, and that vertex 8 has color *green*. This implies that vertex 7 has color *green*, vertices 9 and 5 have color *blue* and vertex 10 has color *red*; but then, $\mathbf{c}(c) = green \neq \mathbf{c}(b)$. This proves the first point of the claim. Observe that each coloring of b and c with distinct colors can be extended to a 3-coloring of H' and H .

As for the second point, since vertices 4 and 6 have color *red*, both 1 and 2 must have color *green*, and since vertex 8 has color *green*, vertex 3 must have color *red*. Consequently, $\mathbf{c}(a) = blue = \mathbf{c}(b)$. ◊

We finally conclude the proof of Theorem 5. By Claim 1, if \mathbf{c} is a 3-coloring of G' then, for any path P_{uv} of \mathcal{P} , we have $\mathbf{c}(u_1) \neq \mathbf{c}(u_2)$ and $\mathbf{c}(u_2) = \mathbf{c}(u_i)$ for all $3 \leq i \leq k$. Hence, \mathbf{c} induces a 3-coloring of G . Conversely, it is easy to see that any 3-coloring of G can be extended to a 3-coloring of G' . □



Fig. 7. The gadget H (left) and a 0-bend CPG representation of it (right).

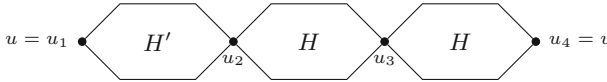


Fig. 8. Adding gadgets H and H' .

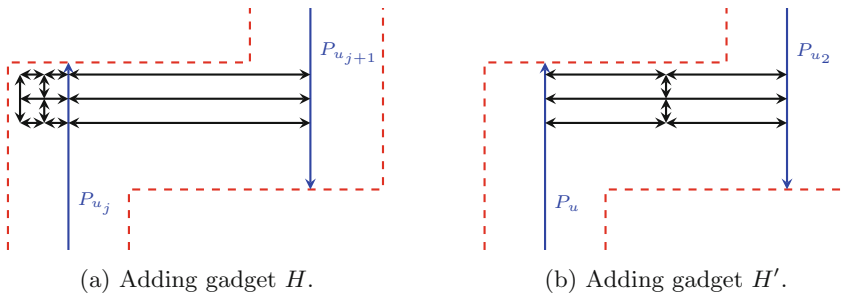


Fig. 9. Locally adding gadgets to control the color of the vertices.

6 Conclusion

We conclude by stating the following open questions:

1. Are B_1 -CPG graphs 5-colorable?
2. Can we characterize those planar graphs which are CPG?
3. Is RECOGNITION NP-complete for B_k -CPG graphs with $k > 0$?

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